

# Dimensions of Points in Self-Similar Fractals

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## Abstract

We use nontrivial connections between the theory of computing and the fine-scale geometry of Euclidean space to give a complete analysis of the dimensions of individual points in fractals that are computably self-similar.

## 1 Introduction

This paper analyzes the dimensions of points in the most widely known type of fractals, the self-similar fractals. Our analysis uses nontrivial connections between the theory of computing and the fine-scale geometry of Euclidean space. In order to explain our results, we briefly review self-similar fractals and the dimensions of points.

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## 1.1 Self-Similar Fractals

The class of self-similar fractals includes such famous objects as the Sierpinski triangle, the Cantor set, the von Koch curve, and the Menger sponge, along with many more exotic sets in Euclidean space [2, 9, 10, 12]. To be concrete, consider the Sierpinski triangle, which is constructed by the process illustrated in Figure 1. We start (at the left) with the equilateral triangle  $D$  whose vertices are the points  $v_0 = (0, 0)$ ,  $v_1 = (1, 0)$ , and  $v_2 = (\frac{1}{2}, \frac{\sqrt{3}}{2})$  in  $\mathbb{R}^2$  (together with this triangle's interior). The construction is carried out by three functions  $S_0, S_1, S_2 : \mathbb{R}^2 \rightarrow \mathbb{R}^2$  defined by

$$S_i(x) = v_i + \frac{1}{2}(x - v_i)$$

for each  $x \in \mathbb{R}^2$  and  $i = 0, 1, 2$ . Note that  $|S_i(x) - S_i(y)| = \frac{1}{2}|x - y|$  always holds, i.e., each  $S_i$  is a *contracting similarity* with *contraction ratio*  $c_i = \frac{1}{2}$ . Note also that each  $S_i$  maps the triangle  $D$  onto a similar subtriangle containing the vertex  $v_i$ .

We use the alphabet  $\Sigma = \{0, 1, 2\}$  to specify the contracting similarities  $S_0, S_1, S_2$ . Each infinite sequence  $T \in \Sigma^\infty$  over this alphabet *codes* a point  $S(T)$  in the Sierpinski triangle via the following recursion. (See Figure 1.) We start at time  $t = 0$  in the triangle  $\Delta_0 = D$ . At time  $t + 1$ , we move into the subtriangle  $\Delta_{t+1}$  of  $\Delta_t$  given by the (appropriately rescaled) contracting similarity  $S_{T[t]}$ , where  $T[t]$  is the  $t^{\text{th}}$  symbol in  $T$ . The point  $S(T)$  is then the unique

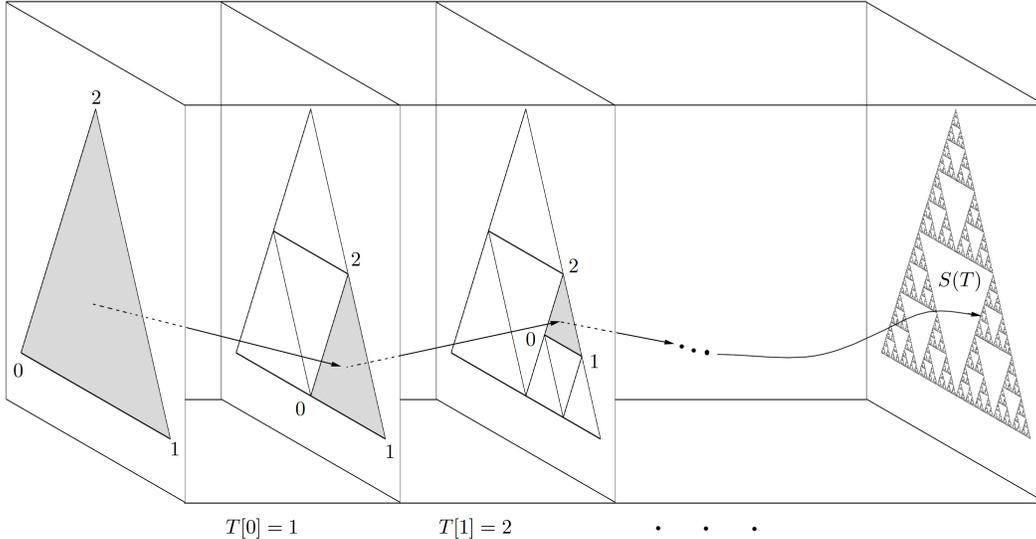


Figure 1. A sequence  $T \in \{0, 1, 2\}^\infty$  codes a point  $S(T)$  in the Sierpinski triangle.

point in  $\mathbb{R}^2$  lying in all the triangles  $\Delta_0, \Delta_1, \Delta_2, \dots$ . Finally, the Sierpinski triangle is the set

$$F(S) = \{S(T) \mid T \in \Sigma_k^\infty\}$$

of all points coded in this fashion.

Self-similar fractals are defined by generalizing the above construction. We work in a Euclidean space  $\mathbb{R}^n$ . An *iterated function system (IFS)* is a list  $S = (S_0, \dots, S_{k-1})$  of two or more contracting similarities  $S_i : \mathbb{R}^n \rightarrow \mathbb{R}^n$  that map an initial nonempty, closed set  $D \subseteq \mathbb{R}^n$  into itself. Each  $S_i$  has a contraction ratio  $c_i \in (0, 1)$ . (The contraction ratios  $c_0, \dots, c_{k-1}$  need not be the same.) The alphabet  $\Sigma_k = \{0, \dots, k-1\}$  is used to specify the contracting similarities in  $S$ , and each sequence  $T \in \Sigma_k^\infty$  codes a point  $S(T) \in \mathbb{R}^n$  in the now-obvious manner. The *attractor* of the IFS  $S$  is the set

$$F(S) = \{S(T) \mid T \in \Sigma_k^\infty\}.$$

In general, the sets  $S_0(D), \dots, S_{k-1}(D)$  may not be disjoint, so a point  $x \in F(S)$  may have many *coding sequences*, i.e.,

many sequences  $T$  for which  $S(T) = x$ . A *self-similar fractal* is a set  $F \subseteq \mathbb{R}^n$  that is the attractor of an IFS  $S$  that satisfies a technical *open set condition* (defined in section 3), which ensures that the sets  $S_0(D), \dots, S_{k-1}(D)$  are “nearly” disjoint.

The *similarity dimension* of a self-similar fractal  $F$  is the (unique) solution  $\text{sdim}(F)$  of the equation

$$\sum_{i=0}^{k-1} c_i^{\text{sdim}(F)} = 1, \quad (1.1)$$

where  $c_0, \dots, c_{k-1}$  are the contraction ratios of any IFS  $S$  satisfying the open set condition and  $F(S) = F$ . A classical theorem of Moran [33] and Falconer [11] says that, for any self-similar fractal  $F$ ,

$$\dim_{\text{H}}(F) = \text{Dim}_{\text{P}}(F) = \text{sdim}(F), \quad (1.2)$$

i.e., the Hausdorff and packing dimensions of  $F$  coincide with its similarity dimension. In addition to its theoretical interest, the Moran-Falconer theorem has the pragmatic consequence that the Hausdorff and packing dimensions of a self-similar fractal are easily computed from

the contraction ratios by solving equation (1.1).

## 1.2 Dimensions of Points

The theory of computing has recently been used to provide a meaningful notion of the *dimensions of individual points* in Euclidean space [29, 1, 17, 30]. These dimensions are robust in that they have many equivalent characterizations. For the purposes of this paper, we define these dimensions in terms of Kolmogorov complexities of rational approximations in Euclidean space.

For each  $x \in \mathbb{R}^n$  and  $r \in \mathbb{N}$ , we define the *Kolmogorov complexity* of  $x$  at *precision*  $r$  to be the natural number

$$K_r(x) = \min\{K(q) \mid q \in \mathbb{Q}^n \text{ and } |q - x| \leq 2^{-r}\},$$

where  $K(q)$  is the Kolmogorov complexity of the rational point  $q$  [25]. That is,  $K_r(x)$  is the minimum length of any program  $\pi \in \{0, 1\}^*$  for which  $U(\pi)$  – the output of a fixed universal Turing machine on input  $\pi$  – is a rational approximation of  $x$  to within  $2^{-r}$ . (Related notions of approximate Kolmogorov complexity have recently been considered by Vitanyi and Vereshchagin [42] and Fortnow, Lee and Vereshchagin [15].)

**Definition.** Let  $x \in \mathbb{R}^n$ .

1. The *dimension* of the point  $x$  is

$$\dim(x) = \liminf_{r \rightarrow \infty} \frac{K_r(x)}{r}. \quad (1.3)$$

2. The *strong dimension* of the point  $x$  is

$$\text{Dim}(x) = \limsup_{r \rightarrow \infty} \frac{K_r(x)}{r}. \quad (1.4)$$

Intuitively,  $\dim(x)$  and  $\text{Dim}(x)$  are the lower and upper asymptotic *information densities* of the point  $x \in \mathbb{R}^n$ .

It is easy to see that  $0 \leq \dim(x) \leq \text{Dim}(x) \leq n$  for all  $x \in \mathbb{R}^n$ . In fact, this is the only restriction that holds in general, i.e., for any two real numbers  $0 \leq \alpha \leq \beta \leq n$ , there is a point  $x \in \mathbb{R}^n$  with  $\dim(x) = \alpha$  and  $\text{Dim}(x) = \beta$  [1]. Points  $x$  that are computable have  $\dim(x) = \text{Dim}(x) = 0$ , while points  $x$  that are random (in the sense of Martin-Löf [31]) have  $\dim(x) = \text{Dim}(x) = n$ .

The dimensions  $\dim(x)$  and  $\text{Dim}(x)$  are well defined and robust, but are they *geometrically* meaningful? Prior work already indicates an affirmative answer. By Hitchcock’s correspondence principle for constructive dimension ([20], extending a result of [39]), together with the absolute stability of constructive dimension [29], if  $X \subseteq \mathbb{R}^n$  is any countable (not necessarily effective) union of computably closed, i.e.,  $\Pi_1^0$ , sets, then

$$\dim_{\text{H}}(X) = \sup_{x \in X} \dim(x). \quad (1.5)$$

That is, the *classical* Hausdorff dimension [12] of any such set is completely determined by the dimensions of its individual points. Many, perhaps most, of the sets which arise in “standard” mathematical practice are unions of computably closed sets, so (1.5) constitutes strong *prima facie* evidence that the dimensions of individual points are indeed geometrically meaningful.

Appendix B shows that the definitions (1.3) and (1.4) are equivalent to the original definitions of dimension and strong dimension [29, 1] and are thus constructive versions of the two most important classical fractal dimensions, namely Hausdorff dimension and packing dimension, respectively.

### 1.3 Our Results

Our main theorem concerns the dimensions of points in fractals that are *computably self-similar*, meaning that they are attractors of *computable* iterated function systems satisfying the open set condition. (We note that most self-similar fractals occurring in practice – including the four famous examples mentioned in section 1.1 – are, in fact, computably self-similar.) Our main theorem says that, if  $F$  is any fractal that is computably self-similar with the IFS  $S$  as witness, then, for every point  $x \in F$  and every coding sequence  $T$  for  $x$ , the dimension and strong dimension of the point  $x$  are given by the dimension formulas

$$\dim(x) = \text{sdim}(F)\dim^{\pi_S}(T) \quad (1.6)$$

and

$$\text{Dim}(x) = \text{sdim}(F)\text{Dim}^{\pi_S}(T), \quad (1.7)$$

where  $\dim^{\pi_S}(T)$  and  $\text{Dim}^{\pi_S}(T)$  are the dimension and strong dimension of  $T$  with respect to the probability measure  $\pi_S$  on the alphabet  $\Sigma_k$  defined by

$$\pi_S(i) = c_i^{\text{sdim}(F)} \quad (1.8)$$

for all  $i \in \Sigma_k$ . (We define  $\dim^{\pi_S}(T)$  and  $\text{Dim}^{\pi_S}(T)$  in the next paragraph.) This theorem gives a complete analysis of the dimensions of points in computably self-similar fractals and the manner in which the dimensions of these points arise from the dimensions of their coding sequences.

In order to understand the right-hand sides of equations (1.6) and (1.7), we now define the dimensions  $\dim^{\pi_S}(T)$  and  $\text{Dim}^{\pi_S}(T)$ .

**Definition.** Let  $\Sigma$  be an alphabet with  $2 \leq |\Sigma| < \infty$ , and let  $\pi$  be a positive probability measure on  $\Sigma$ . Let  $w \in \Sigma^*$  and  $T \in \Sigma^\infty$ .

1. The Shannon *self-information* of  $w$  with respect to  $\pi$  is

$$\mathcal{I}_\pi(w) = \log \frac{1}{\pi(w)} = \sum_{i=0}^{|w|-1} \log \frac{1}{\pi(w[i])}, \quad (1.9)$$

where the logarithm is base-2 [8].

2. The *dimension* of  $T$  with respect to  $\pi$  is

$$\dim^\pi(T) = \liminf_{j \rightarrow \infty} \frac{K(T[0..j-1])}{\mathcal{I}_\pi(T[0..j-1])}. \quad (1.10)$$

3. The *strong dimension* of  $T$  with respect to  $\pi$  is

$$\text{Dim}^\pi(T) = \limsup_{j \rightarrow \infty} \frac{K(T[0..j-1])}{\mathcal{I}_\pi(T[0..j-1])}. \quad (1.11)$$

The dimensions  $\dim^\pi(T)$  and  $\text{Dim}^\pi(T)$  are measures of the algorithmic information density of  $T$ , but the “density” here is now an information-to-cost ratio. In this ratio, the “information” is algorithmic information, i.e., Kolmogorov complexity, and the “cost” is the Shannon self-information. To see why this makes sense, consider the case of interest in our main theorem. In this case, (1.8) says that the cost of a string  $w \in \Sigma_k^*$  is

$$\mathcal{I}_\pi(w) = \text{sdim}(F) \sum_{j=0}^{|w|-1} \log \frac{1}{c_{w[j]}},$$

i.e., the sum of the costs of the symbols in  $w$ , where the cost of a symbol  $i \in \Sigma_k$  is  $\text{sdim}(F) \log(1/c_i)$ . These symbol costs are computational and realistic. A symbol  $i$  with high cost invokes a similarity  $S_i$  with a small contraction ratio  $c_i$ , thereby necessitating a high-precision computation.

Appendix A shows that definitions (1.10) and (1.11) are equivalent to “gale

characterizations” of these dimensions, and hence that  $\dim^\pi(T)$  is a constructive version of Billingsley dimension [3, 7].

Although our main theorem only applies directly to computably self-similar fractals, we use relativization to show that the Moran-Falconer theorem (1.2) for arbitrary self-similar fractals is an easy consequence of our main theorem. Hence, as is often the case, a theorem of computable analysis (i.e., the theoretical foundations of scientific computing [5]) has an immediate corollary in classical analysis.

The proof of our main theorem has some geometric and combinatorial similarities with the classical proofs of Moran [33] and Falconer [11], but the argument here is information-theoretic. As such, it gives a more clear understanding of the computational aspects of dimension in self-similar fractals, even in the classical case.

We note that Cai and Hartmanis [6] and Fernau and Staiger [14] have conducted related investigations of Hausdorff dimension in iterated function fractals and their coding spaces, but with different motivations and results. Our focus here is on a *pointwise* analysis of dimensions.

Some of the most difficult open problems in geometric measure theory involve establishing lower bounds on the fractal dimensions of various sets. Kolmogorov complexity has proven to be a powerful tool for lower-bound arguments, leading to the solution of many long-standing open problems in discrete mathematics [25]. There is thus reason to hope that our pointwise approach to fractal dimension, coupled with the introduction of Kolmogorov complexity techniques, will lead to progress in this classical area. In any case, our results extend computable analysis [34, 22, 44] in a new, geometric direction.

## 2 Preliminaries

Given a finite alphabet  $\Sigma$ , we write  $\Sigma^*$  for the set of all (finite) *strings* over  $\Sigma$  and  $\Sigma^\infty$  for the set of all (infinite) *sequences* over  $\Sigma$ . If  $\psi \in \Sigma^* \cup \Sigma^\infty$  and  $0 \leq i \leq j < |\psi|$ , where  $|\psi|$  is the length of  $\psi$ , then  $\psi[i]$  is the  $i$ th symbol in  $\psi$  (where  $\psi[0]$  is the leftmost symbol in  $\psi$ ), and  $\psi[i..j]$  is the string consisting of the  $i$ th through the  $j$ th symbols in  $\psi$ . If  $w \in \Sigma^*$  and  $\psi \in \Sigma^* \cup \Sigma^\infty$ , then  $w$  is a prefix of  $\psi$ , and we write  $w \sqsubseteq \psi$ , if there exists  $i \in \mathbb{N}$  such that  $w = \psi[0..i-1]$ .

For functions on Euclidean space, we use the computability notion formulated by Grzegorzczuk [16] and Lacombe [23] in the 1950’s and expounded in the monographs by Pour-El and Richards [34], Ko [22], and Weihrauch [44] and in the recent survey paper by Braverman and Cook [5]. A function  $f : \mathbb{R}^n \rightarrow \mathbb{R}^n$  is *computable* if there is an oracle Turing machine  $M$  with the following property. For all  $x \in \mathbb{R}^n$  and  $r \in \mathbb{N}$ , if  $M$  is given a function oracle  $\varphi_x : \mathbb{N} \rightarrow \mathbb{Q}^n$  such that, for all  $k \in \mathbb{N}$ ,  $|\varphi_x(k) - x| \leq 2^{-k}$ , then  $M$ , with oracle  $\varphi_x$  and input  $r$ , outputs a rational point  $M^{\varphi_x}(r) \in \mathbb{Q}^n$  such that  $|M^{\varphi_x}(r) - f(x)| \leq 2^{-r}$ .

A point  $x \in \mathbb{R}^n$  is *computable* if there is a computable function  $\psi_x : \mathbb{N} \rightarrow \mathbb{Q}^n$  such that, for all  $r \in \mathbb{N}$ ,  $|\psi_x(r) - x| \leq 2^{-r}$ .

For subsets of Euclidean space, we use the computability notion introduced by Brattka and Weihrauch [4] (see also [44, 5]). A set  $X \subseteq \mathbb{R}^n$  is *computable* if there is a computable function  $f_X : \mathbb{Q}^n \times \mathbb{N} \rightarrow \{0, 1\}$  that satisfies the following two conditions for all  $q \in \mathbb{Q}^n$  and  $r \in \mathbb{N}$ .

- (i) If there exists  $x \in X$  such that  $|x - q| \leq 2^{-r}$ , then  $f_X(q, r) = 1$ .
- (ii) If there is no  $x \in X$  such that  $|x - q| \leq 2^{1-r}$ , then  $f_X(q, r) = 0$ .

The following two observations are well known and easy to verify.

**Observation 2.1** *A set  $X \subseteq \mathbb{R}^n$  is computable if and only if the associated distance function*

$$\begin{aligned} \rho_X : \mathbb{R}^n &\rightarrow [0, \infty) \\ \rho_X(y) &= \inf_{x \in X} |x - y| \end{aligned}$$

*is computable.*

**Observation 2.2** *If  $X \subseteq \mathbb{R}^n$  is both computable and closed, then  $X$  is a computably closed, i.e.,  $\Pi_1^0$ , set.*

All logarithms in this paper are base-2.

### 3 More on Self-Similar Fractals

This expository section reviews a fragment of the theory of self-similar fractals that is adequate for understanding our main theorem and its proof. Our treatment is self-contained, but of course far from complete. The interested reader is referred to any of the standard texts [2, 9, 10, 12] for more extensive discussion.

**Definition.** An *iterated function system (IFS)* is a finite sequence  $S = (S_0, \dots, S_{k-1})$  of two or more contracting similarities on a nonempty, closed set  $D \subseteq \mathbb{R}^n$ . We call  $D$  the *domain* of  $S$ , writing  $D = \text{dom}(S)$ .

We use the standard notation  $\mathcal{K}(D)$  for the set of all nonempty compact (i.e., closed and bounded) subsets of a nonempty closed set  $D \subseteq \mathbb{R}^n$ . For each IFS  $S$ , we write  $\mathcal{K}(S) = \mathcal{K}(\text{dom}(S))$ .

For each IFS  $S = (S_0, \dots, S_{k-1})$ , we define the transformation  $S : \mathcal{K}(S) \rightarrow \mathcal{K}(S)$  by

$$S(A) = \bigcup_{i=0}^{k-1} S_i(A)$$

for all  $A \in \mathcal{K}(S)$ , where  $S_i(A)$  is the image of  $A$  under the contracting similarity  $S_i$ .

**Observation 3.1** *For each IFS  $S$ , there exists  $A \in \mathcal{K}(S)$  such that  $S(A) \subseteq A$ .*

For each IFS  $S = (S_0, \dots, S_{k-1})$  and each set  $A \in \mathcal{K}(S)$  satisfying  $S(A) \subseteq A$ , we define the function  $S_A : \Sigma_k^* \rightarrow \mathcal{K}(S)$  by the recursion

$$S_A(\lambda) = A;$$

$$S_A(iw) = S_i(S_A(w))$$

for all  $w \in \Sigma_k^*$  and  $i \in \Sigma_k$ .

If  $c = \max\{c_0, \dots, c_{k-1}\}$ , where  $c_0, \dots, c_{k-1}$  are contraction ratios of  $S_0, \dots, S_{k-1}$ , respectively, then routine inductions establish that, for all  $w \in \Sigma_k^*$  and  $i \in \Sigma_k$ ,

$$S_A(iw) \subseteq S_A(w) \quad (3.1)$$

and

$$\text{diam}(S_A(w)) \leq c^{|w|} \text{diam}(A). \quad (3.2)$$

Since  $c \in (0, 1)$ , it follows that, for each sequence  $T \in \Sigma_k^\infty$ , there is a unique point  $S_A(T) \in \mathbb{R}^n$  such that

$$\bigcap_{w \sqsubseteq T} S_A(w) = \{S_A(T)\}. \quad (3.3)$$

In this manner, we have defined a function  $S_A : \Sigma_k^\infty \rightarrow \mathbb{R}^n$ . The following observation shows that this function does not really depend on the choice of  $A$ .

**Observation 3.2** *Let  $S$  be an IFS. If  $A, B \in \mathcal{K}(S)$  satisfy  $S(A) \subseteq A$  and  $S(B) \subseteq B$ , then  $S_A = S_B$ .*

For each IFS  $S$ , we define the induced function  $S : \Sigma_k^\infty \rightarrow \mathbb{R}^n$  by setting  $S = S_A$ , where  $A$  is any element of  $\mathcal{K}(S)$  satisfying  $S(A) \subseteq A$ . By Observations 3.1 and 3.2, this induced function  $S$  is well-defined.

We now have the machinery to define a rich collection of fractals in  $\mathbb{R}^n$ .

**Definition.** The *attractor* (or *invariant set*) of an IFS  $S = (S_0, \dots, S_{k-1})$  is the set

$$F(S) = S(\Sigma_k^\infty),$$

i.e., the range of the induced function  $S : \Sigma_k^\infty \rightarrow \mathbb{R}^n$ .

It is well-known that the attractor  $F(S)$  is the unique fixed point of the induced transformation  $S : \mathcal{K}(S) \rightarrow \mathcal{K}(S)$ , but we do not use this fact here.

For each  $T \in \Sigma_k^\infty$ , we call  $T$  a *coding sequence*, or an *S-code*, of the point  $S(T) \in F(S)$ .

In general, the attractor of an IFS  $S = (S_0, \dots, S_{k-1})$  is easiest to analyze when the sets  $S_0(\dim(S)), \dots, S_{k-1}(\dim(S))$  are “nearly disjoint”. (Intuitively, this prevents each point  $x \in F(S)$  from having “too many” coding sequences  $T \in \Sigma_k^\infty$ .) The following definition makes this notion precise.

**Definition.** An IFS  $S = (S_0, \dots, S_{k-1})$  with domain  $D$  satisfies the *open set condition* if there exists a nonempty, bounded, open set  $G \subseteq D$  such that  $S_0(G), \dots, S_{k-1}(G)$  are disjoint subsets of  $G$ .

We now define the most widely known type of fractal.

**Definition.** A *self-similar fractal* is a set  $F \subseteq \mathbb{R}^n$  that is the attractor of an IFS that satisfies the open set condition.

## 4 Pointwise Analysis of Dimensions

In this section we prove our main theorem, which gives a precise analysis of the dimensions of individual points in computably self-similar fractals. We first recall the known fact that such fractals are

computable.

**Definition.** An IFS  $S = (S_0, \dots, S_{k-1})$  is computable if  $\text{dom}(S)$  is a computable set and the functions  $S_0, \dots, S_{k-1}$  are computable.

**Theorem 4.1** (*Kamo and Kawamura [21]*). *For every computable IFS  $S$ , the attractor  $F(S)$  is a computable set.*

One consequence of Theorem 4.1 is the following.

**Corollary 4.2** *For every computable IFS  $S$ ,  $\text{cdim}(F(S)) = \text{dim}_H(F(S))$ .*

We next present three lemmas that we use in the proof of our main theorem. The first is a well-known geometric fact (e.g., it is Lemma 9.2 in [12]).

**Lemma 4.3** *Let  $\mathcal{G}$  be a collection of disjoint open sets in  $\mathbb{R}^n$ , and let  $r, a, b \in (0, \infty)$ . If every element of  $\mathcal{G}$  contains a ball of radius  $ar$  and is contained in a ball of radius  $br$ , then no ball of radius  $r$  meets more than  $(\frac{1+2b}{a})^n$  of the closures of the elements of  $\mathcal{G}$ .*

Our second lemma gives a computable means of assigning rational “hubs” to the various open sets arising from a computable IFS satisfying the open set condition.

**Definition.** A *hub function* for an IFS  $S = (S_0, \dots, S_{k-1})$  satisfying the open set condition with  $G$  as witness is a function  $h : \Sigma_k^* \rightarrow \mathbb{R}^n$  such that  $h(w) \in S_G(w)$  for all  $w \in \Sigma_k^*$ . In this case, we call  $h(w)$  the *hub* that  $h$  assigns to the set  $S_G(w)$ .

**Lemma 4.4** *If  $S = (S_0, \dots, S_{k-1})$  is a computable IFS satisfying the open set condition with  $G$  as witness, then there is an exactly computable, rational-valued hub function  $h : \Sigma_k^* \rightarrow \mathbb{Q}^n$  for  $S$  and  $G$ .*

For  $w \in \Sigma_k^*$ , we use the abbreviation  $\mathcal{I}_S(w) = \mathcal{I}_{\pi_S}(w)$ , where  $\pi_S$  is the probability measure defined in section 1.3.

Our third lemma provides a decidable set of well-behaved ‘‘canonical prefixes’’ of sequences in  $\Sigma_k^\infty$ .

**Lemma 4.5** *Let  $S = (S_0, \dots, S_{k-1})$  be a computable IFS, and let  $c_{\min}$  be the minimum of the contraction ratios of  $S = (S_0, \dots, S_{k-1})$ . For any real number*

$$\alpha > \text{sdim}(S) \log \frac{1}{c_{\min}}, \quad (4.1)$$

*there exists a decidable set  $A \subseteq \mathbb{N} \times \Sigma_k^*$  such that, for each  $r \in \mathbb{N}$ , the set*

$$A_r = \{w \in \Sigma_k^* \mid (r, w) \in A\}$$

*has the following three properties.*

- (i) *No element of  $A_r$  is a proper prefix of any element of  $A_{r'}$  for any  $r' \leq r$ .*
- (ii) *Each sequence in  $\Sigma_k^\infty$  has a (unique) prefix in  $A_r$ .*
- (iii) *For all  $w \in A_r$ ,*

$$r \text{sdim}(S) < \mathcal{I}_S(w) < r \text{sdim}(S) + \alpha. \quad (4.2)$$

Our main theorem concerns the following type of fractal.

**Definition.** A *computably self-similar fractal* is a set  $F \subseteq \mathbb{R}^n$  that is the attractor of an IFS that is computable and satisfies the open set condition.

Most self-similar fractals occurring in the literature are, in fact, computably self-similar.

We now have the machinery to give a complete analysis of the dimensions of points in computably self-similar fractals.

**Theorem 4.6 (main theorem).** *If  $F \subseteq \mathbb{R}^n$  is a computably self-similar fractal and  $S$  is an IFS testifying this fact, then, for all points  $x \in F$  and all  $S$ -codes  $T$  of  $x$ ,*

$$\dim(x) = \text{sdim}(F) \dim^{\pi_S}(T) \quad (4.3)$$

and

$$\text{Dim}(x) = \text{sdim}(F) \text{Dim}^{\pi_S}(T). \quad (4.4)$$

**Proof.** Assume the hypothesis, with  $S = (S_0, \dots, S_{k-1})$ . Let  $c_0, \dots, c_{k-1}$  be the contraction ratios of  $S_0, \dots, S_{k-1}$ , respectively, and let  $G$  be a witness to the fact that  $S$  satisfies the open set condition, and let

$l = \max\{0, \lceil \log \text{diam}(G) \rceil\}$ . Let  $h : \Sigma_k^* \rightarrow \mathbb{Q}^n$  be an exactly computable, rational-valued hub function for  $S$  and  $G$  as given by Lemma 4.4. Let  $\alpha = 1 + \text{sdim}(F) \log \frac{1}{c_{\min}}$ , for  $c_{\min} = \min\{c_0, \dots, c_{k-1}\}$ , and choose a decidable set  $A$  for  $S$  and  $\alpha$  as in Lemma 4.5.

For all  $w \in \Sigma_k^*$ , we have

$$\begin{aligned} \text{diam}(S_G(w)) &= \text{diam}(G) \prod_{i=0}^{|w|-1} c_{w[i]} \\ &= \text{diam}(G) \pi_S(w)^{\frac{1}{\text{sdim}(F)}}. \end{aligned}$$

It follows by (4.2) that, for all  $r \in \mathbb{N}$  and  $w \in A_{r+l}$ ,

$$2^{-r} a_1 \leq \text{diam}(S_G(w)) \leq 2^{-r}, \quad (4.5)$$

where  $a_1 = 2^{-\frac{l+\alpha}{\text{sdim}(F)}} \text{diam}(G)$ .

Let  $x \in F$ , and let  $T \in \Sigma_k^\infty$  be an  $S$ -code of  $x$ , i.e.,  $S(T) = x$ . For each  $r \in \mathbb{N}$ , let  $w_r$  be the unique element of  $A_{r+l}$  that is a prefix of  $T$ . Much of this proof is devoted to deriving a close relationship between the Kolmogorov complexities  $K_r(x)$  and  $K(w_r)$ . Once we have this relationship, we will use it to prove (4.3) and (4.4).

Since the hub function  $h$  is computable, there is a constant  $a_2$  such that, for all  $w \in \Sigma_k^*$ ,

$$K(h(w)) \leq K(w) + a_2. \quad (4.6)$$

Since  $h(w_r) \in S_G(w_r)$  and  $x = S(T) \in S_{\overline{G}}(w_r) = \overline{S_G(w_r)}$ , (4.5) tells us that

$$|h(w_r) - x| \leq \text{diam}(S_G(w_r)) \leq 2^{-r},$$

whence

$$K_r(x) \leq K(h(w_r))$$

for all  $r \in \mathbb{N}$ . It follows by (4.6) that

$$K_r(x) \leq K(w_r) + a_2 \quad (4.7)$$

for all  $r \in \mathbb{N}$ . Combining (4.7) and the right-hand inequality in (4.2) gives

$$\frac{K_r(x)}{\text{rsdim}(F)} \leq \frac{K(w_r) + a_2}{\mathcal{I}_S(w_r) - \alpha} \quad (4.8)$$

for all  $r \in \mathbb{N}$ .

Let  $E$  be the set of all triples  $(q, r, w)$  such that  $q \in \mathbb{Q}^n$ ,  $r \in \mathbb{N}$ ,  $w \in A_{r+l}$ , and

$$|q - h(w)| \leq 2^{1-r}. \quad (4.9)$$

Since the set  $A$  and the condition (4.9) are decidable, the set  $E$  is decidable.

For each  $q \in \mathbb{Q}^n$  and  $r \in \mathbb{N}$ , let

$$E_{q,r} = \{w \in \Sigma_k^* \mid (q, r, w) \in E\}.$$

We prove two key properties of the sets  $E_{q,r}$ . First, for all  $q \in \mathbb{Q}^n$  and  $r \in \mathbb{N}$ ,

$$|q - x| \leq 2^{-r} \Rightarrow w_r \in E_{q,r}. \quad (4.10)$$

To see that this holds, assume that  $|q - x| \leq 2^{-r}$ . Since  $x = S(T) \in S_{\overline{G}}(w_r) = \overline{S_G(w_r)}$ , the right-hand inequality in (4.5) tells us that

$$|q - h(w_r)| \leq |q - x| + |x - h(w_r)|$$

$$\leq 2^{-r} + \text{diam}(S_G(w_r)) \leq 2^{1-r},$$

confirming (4.10).

The second key property of the sets  $E_{q,r}$  is that they are small, namely, that

$$|E_{q,r}| \leq \gamma \quad (4.11)$$

holds for all  $q \in \mathbb{Q}^n$  and  $r \in \mathbb{N}$ , where  $\gamma$  is a constant that does not depend on  $q$  or  $r$ . To see this, let  $w \in E_{q,r}$ . Then  $w \in A_{r+l}$  and  $|q - h(w)| \leq 2^{1-r}$ , so  $h(w) \in S_G(w) \cap B(q, 2^{1-r})$ . This argument establishes that

$$w \in E_{q,r} \Rightarrow B(q, 2^{1-r}) \text{ meets } S_G(w). \quad (4.12)$$

Now let

$$\mathcal{G}_r = \{S_G(w) \mid w \in A_{r+l}\}.$$

By our choice of  $G$ ,  $\mathcal{G}_r$  is a collection of disjoint open sets in  $\mathbb{R}^n$ . By the right-hand inequality in (4.5), each element of  $\mathcal{G}_r$  is contained in a closed ball of radius  $2^{-r}$ . Since  $G$  is open, it contains a closed ball of some radius  $a_3 > 0$ . It follows by the left-hand inequality in (4.5) that  $S_G(w)$ , being a contraction of  $G$ , contains a closed ball of radius  $2^{1-r}a_4$ , where  $a_4 = \frac{a_1 a_3}{2 \text{diam}(G)}$ . By Lemma 4.3, this implies that  $B(q, 2^{1-r})$  meets no more than  $\gamma$  of the (closures of the) elements of  $\mathcal{G}_r$ , where  $\gamma = \left(\frac{2}{a_4}\right)^n$ . By (4.12), this confirms (4.11).

Now let  $M$  be a prefix Turing machine with the following property. If  $U(\pi) = q \in \mathbb{Q}^n$  (where  $U$  is the universal prefix Turing machine),  $s_r$  is the  $r^{\text{th}}$  string in a standard enumeration  $s_0, s_1, \dots$  of  $\{0, 1\}^*$ , and  $0 \leq m < |E_{q,r}|$ , then  $M(\pi 0^{|s_r|} 1 s_r 0^m 1)$  is the  $m^{\text{th}}$  element of  $E_{q,r}$ . There is a constant  $a_5$  such that, for all  $w \in \Sigma_k^*$ ,

$$K(w) \leq K_M(w) + a_5. \quad (4.13)$$

Taking  $\pi$  to be a program testifying to the value of  $K_r(x)$  and applying (4.10) and

(4.11) shows that

$$\begin{aligned} K_M(w_r) &\leq K_r(x) + 2\log(r+1) + |E_{q,r}| + 1 \\ &\leq K_r(x) + 2\log(r+1) + \gamma + 1, \end{aligned}$$

whence (4.13) tells us that

$$K(w_r) \leq K_r(x) + \epsilon(r) \quad (4.14)$$

for all  $r \in \mathbb{N}$ , where  $\epsilon(r) = 2\log(r+1) + a_5 + \gamma + 1$ . Combining (4.14) and the right-hand inequality in (4.2) gives

$$\frac{K_r(x)}{r\text{sdim}(F)} \geq \frac{K(w_r) - \epsilon(r)}{\mathcal{I}_S(w_r)} \quad (4.15)$$

for all  $r \in \mathbb{N}$ . Note that  $\epsilon(r) = o(\mathcal{I}_S(w_r))$  as  $r \rightarrow \infty$ .

By (4.8) and (4.15), we now have

$$\frac{K(w_r) - \epsilon(r)}{\mathcal{I}_S(w_r) - \alpha} \leq \frac{K_r(x)}{r\text{sdim}(F)} \leq \frac{K(w_r) + a_1}{\mathcal{I}_S(w_r) - \beta} \quad (4.16)$$

for all  $r \in \mathbb{N}$ . In order to use this relationship between  $K_r(x)$  and  $K(w_r)$ , we need to know that the asymptotic behavior of  $\frac{K(w_r)}{\mathcal{I}_S(w_r)}$  for  $r \in \mathbb{N}$  is the same as the asymptotic behavior of  $\frac{K(w)}{\mathcal{I}_S(w)}$  for arbitrary prefixes  $w$  of  $T$ . Our verification of this fact makes repeated use of the *additivity* of  $\mathcal{I}_S$ , by which we mean that

$$\mathcal{I}_S(uv) = \mathcal{I}_S(u) + \mathcal{I}_S(v) \quad (4.17)$$

holds for all  $u, v \in \Sigma_k^*$ .

Let  $r \in \mathbb{N}$ , and let  $w_r \sqsubseteq w \sqsubseteq w_{r+1}$ , writing  $w = w_r u$  and  $w_{r+1} = wv$ . Then (4.17) tells us that

$$\mathcal{I}_S(w_r) \leq \mathcal{I}_S(w) \leq \mathcal{I}_S(w_{r+1}),$$

and (4.2) tells us that

$$\mathcal{I}_S(w_{r+1}) - \mathcal{I}_S(w_r) \leq \text{sdim}(F) + \alpha,$$

so we have

$$\mathcal{I}_S(w_r) \leq \mathcal{I}_S(w) \leq \mathcal{I}_S(w_r) + a_6, \quad (4.18)$$

where  $a_6 = \text{sdim}(F) + \alpha$ . We also have

$$\begin{aligned} a_6 &\geq \mathcal{I}_S(w_{r+1}) - \mathcal{I}_S(w_r) \\ &= \mathcal{I}_S(wv) \\ &= \log \frac{1}{\pi_S(wv)} \\ &\geq \log c_{\min}^{-\text{sdim}(F)|uv|} \\ &= |uv| \text{sdim}(F) \log \frac{1}{c_{\min}}, \end{aligned}$$

i.e.,

$$|w_{r+1}| - |w_r| \leq a_7, \quad (4.19)$$

where  $a_7 = \frac{a_6}{\text{sdim}(F) \log \frac{1}{c_{\min}}}$ .

Since (4.19) holds for all  $r \in \mathbb{N}$  and  $a_7$  does not depend on  $r$ , there is a constant  $a_8$  such that, for all  $r \in \mathbb{N}$  and  $w_r \sqsubseteq w \sqsubseteq w_{r+1}$ ,

$$|K(w) - K(w_r)| \leq a_8. \quad (4.20)$$

It follows by (4.18) that

$$\frac{K(w_r) - a_8}{\mathcal{I}_S(w_r) + a_6} \leq \frac{K(w)}{\mathcal{I}_S(w)} \leq \frac{K(w_r) + a_8}{\mathcal{I}_S(w_r)} \quad (4.21)$$

holds for all  $r \in \mathbb{N}$  and  $w_r \sqsubseteq w \sqsubseteq w_{r+1}$ .

By (4.16), (4.21), Theorem B.5, and Theorem B.1, we now have

$$\begin{aligned} \dim(x) &= \liminf_{r \rightarrow \infty} \frac{K_r(x)}{r} \\ &= \text{sdim}(F) \liminf_{r \rightarrow \infty} \frac{K(w_r)}{\mathcal{I}_S(w_r)} \\ &= \text{sdim}(F) \liminf_{j \rightarrow \infty} \frac{K(T[0..j-1])}{\mathcal{I}_S(T[0..j-1])} \\ &= \text{sdim}(F) \dim^{\pi_S}(T) \end{aligned}$$

and, similarly,

$$\text{Dim}(x) = \text{sdim}(F) \text{Dim}^{\pi_S}(T). \quad \square$$

Finally, we use relativization to derive the following well-known classical theorem from our main theorem.

**Corollary 4.7** (*Moran [33], Falconer [11]*). *For every self-similar fractal  $F \subseteq \mathbb{R}^n$ ,*

$$\dim_{\text{H}}(F) = \text{Dim}_{\text{P}}(F) = \text{sdim}(F).$$

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# Technical Appendix

## A Dimensions relative to probability measures

Here we develop the basic theory of constructive fractal dimension on a sequence space  $\Sigma^\infty$  with respect to a suitable probability measure on  $\Sigma^\infty$ . We first review the classical Hausdorff and packing dimensions.

Let  $\rho$  be a metric on a set  $\mathcal{X}$ . We use the following standard terminology. The *diameter* of a set  $X \subseteq \mathcal{X}$  is

$$\text{diam}(X) = \sup \{ \rho(x, y) \mid x, y \in X \}$$

(which may be  $\infty$ ). For each  $x \in \mathcal{X}$  and  $r \in \mathbb{R}$ , the *closed ball* of radius  $r$  about  $x$  is the set

$$B(x, r) = \{ y \in \mathcal{X} \mid \rho(y, x) \leq r \},$$

and the *open ball* of radius  $r$  about  $x$  is the set

$$B^o(x, r) = \{ y \in \mathcal{X} \mid \rho(y, x) < r \}.$$

A *ball* is any set of the form  $B(x, r)$  or  $B^o(x, r)$ . A ball  $B$  is *centered* in a set  $X \subseteq \mathcal{X}$  if  $B = B(x, r)$  or  $B = B^o(x, r)$  for some  $x \in X$  and  $r \geq 0$ .

For each  $\delta > 0$ , we let  $\mathcal{C}_\delta$  be the set of all countable collections  $\mathcal{B}$  of balls such that  $\text{diam}(B) \leq \delta$  for all  $B \in \mathcal{B}$ , and we let  $\mathcal{D}_\delta$  be the set of all  $\mathcal{B} \in \mathcal{C}_\delta$  such that the balls in  $\mathcal{B}$  are pairwise disjoint. For each  $X \subseteq \mathcal{X}$  and  $\delta > 0$ , we define the sets

$$\begin{aligned} \mathcal{H}_\delta(X) &= \left\{ \mathcal{B} \in \mathcal{C}_\delta \mid X \subseteq \bigcup_{B \in \mathcal{B}} B \right\}, \\ \mathcal{P}_\delta(X) &= \{ \mathcal{B} \in \mathcal{D}_\delta \mid (\forall B \in \mathcal{B}) B \text{ is centered in } X \}. \end{aligned}$$

If  $\mathcal{B} \in \mathcal{H}_\delta(X)$ , then we call  $\mathcal{B}$  a  $\delta$ -cover of  $X$ . If  $\mathcal{B} \in \mathcal{P}_\delta(X)$ , then we call  $\mathcal{B}$  a  $\delta$ -packing of  $X$ . For  $X \subseteq \mathcal{X}$ ,  $\delta > 0$  and  $s \geq 0$ , we define the quantities

$$\begin{aligned} H_\delta^s(X) &= \inf_{\mathcal{B} \in \mathcal{H}_\delta(X)} \sum_{B \in \mathcal{B}} \text{diam}(B)^s, \\ P_\delta^s(X) &= \sup_{\mathcal{B} \in \mathcal{P}_\delta(X)} \sum_{B \in \mathcal{B}} \text{diam}(B)^s. \end{aligned}$$

Since  $H_\delta^s(X)$  and  $P_\delta^s(X)$  are monotone as  $\delta \rightarrow 0$ , the limits

$$\begin{aligned} H^s(X) &= \lim_{\delta \rightarrow 0} H_\delta^s(X), \\ P_0^s(X) &= \lim_{\delta \rightarrow 0} P_\delta^s(X) \end{aligned}$$

exist, though they may be infinite. Let

$$P^s(X) = \inf \left\{ \sum_{i=0}^{\infty} P_0^s(X_i) \mid X \subseteq \bigcup_{i=0}^{\infty} X_i \right\}. \quad (\text{A.1})$$

It is routine to verify that the set functions  $H^s$  and  $P^s$  are outer measures [12]. The quantities  $H^s(X)$  and  $P^s(X)$  – which may be infinite – are called the  $s$ -dimensional Hausdorff (outer) ball measure and the  $s$ -dimensional packing (outer) ball measure of  $X$ , respectively. The optimization (A.1) over all countable partitions of  $X$  is needed because the set function  $P_0^s$  is *not* an outer measure.

**Definition.** Let  $\rho$  be a metric on a set  $\mathcal{X}$ , and let  $X \subseteq \mathcal{X}$ .

1. (Hausdorff [18]). The *Hausdorff dimension* of  $X$  with respect to  $\rho$  is

$$\dim_{\mathbb{H}}^{(\rho)}(X) = \inf \{s \in [0, \infty) \mid H^s(X) = 0\}.$$

2. (Tricot [41], Sullivan [40]). The *packing dimension* of  $X$  with respect to  $\rho$  is

$$\text{Dim}_{\mathbb{P}}^{(\rho)}(X) = \inf \{s \in [0, \infty) \mid P^s(X) = 0\}.$$

When  $\mathcal{X}$  is a Euclidean space  $\mathbb{R}^n$  and  $\rho$  is the usual Euclidean metric on  $\mathbb{R}^n$ ,  $\dim_{\mathbb{H}}^{(\rho)}$  and  $\text{Dim}_{\mathbb{P}}^{(\rho)}$  are the ordinary Hausdorff and packing dimensions, also denoted by  $\dim_{\mathbb{H}}$  and  $\text{Dim}_{\mathbb{P}}$ , respectively.

We now focus our attention on sequence spaces. Let  $\Sigma$  be a finite alphabet with  $|\Sigma| \geq 2$ . A (Borel) probability measure on  $\Sigma^\infty$  is a function  $\nu : \Sigma^* \rightarrow [0, 1]$  such that  $\nu(\lambda) = 1$  and  $\nu(w) = \sum_{a \in \Sigma} \nu(wa)$  for all  $w \in \Sigma^*$ . Intuitively,  $\nu(w)$  is the probability that  $w \sqsubseteq S$  when a sequence  $S \in \Sigma^\infty$  is chosen according to the probability measure  $\nu$ . A probability measure  $\nu$  on  $\Sigma^\infty$  is *strongly positive* if there exists  $\delta > 0$  such that, for all  $w \in \Sigma^*$  and  $a \in \Sigma$ ,  $\nu(wa) > \delta\nu(w)$ .

The following type of probability measure is used in our main theorem.

**Example A.1** Let  $\pi$  be a probability measure on the alphabet  $\Sigma$ , i.e., a function  $\pi : \Sigma \rightarrow [0, 1]$  such that  $\sum_{a \in \Sigma} \pi(a) = 1$ . Then  $\pi$  induces the product probability measure  $\pi$  on  $\Sigma^\infty$  defined by

$$\pi(w) = \prod_{i=0}^{|w|-1} \pi(w[i])$$

for all  $w \in \Sigma^*$ . If  $\pi$  is positive on  $\Sigma$ , i.e.,  $\pi(a) > 0$  for all  $a \in \Sigma$ , then the probability measure  $\pi$  on  $\Sigma^\infty$  is strongly positive.

**Example A.2** We reserve the symbol  $\mu$  for the uniform probability measure on  $\Sigma^\infty$ , which is the function  $\mu : \Sigma^* \rightarrow [0, \infty)$  defined by

$$\mu(w) = |\Sigma|^{-|w|}$$

for all  $w \in \Sigma^*$ . Note that this is the special case of Example A.1 in which  $\pi(a) = 1/|\Sigma|$  for each  $a \in \Sigma$ .

**Definition.** The metric induced by a strongly positive probability measure  $\nu$  on  $\Sigma^\infty$  is the function  $\rho_\nu : \Sigma^\infty \times \Sigma^\infty \rightarrow [0, 1]$  defined by

$$\rho_\nu(S, T) = \inf \{\nu(w) \mid w \sqsubseteq S \text{ and } w \sqsubseteq T\}$$

for all  $S, T \in \Sigma^\infty$ .

The following fact is easily verified.

**Observation A.3** For every strongly positive probability measure  $\nu$  on  $\Sigma^\infty$ , the function  $\rho_\nu$  is a metric on  $\Sigma^\infty$ .

Hausdorff and packing dimensions with respect to probability measures on sequence spaces are defined as follows.

**Definition.** Let  $\Sigma$  be a finite alphabet with  $|\Sigma| \geq 2$ , let  $\nu$  be a strongly positive probability measure on  $\Sigma^\infty$ , and let  $X \subseteq \Sigma^\infty$ .

1. The *Hausdorff dimension* of  $X$  with respect to  $\nu$  (also called the *Billingsley dimension* of  $X$  with respect to  $\nu$  [3, 7]) is

$$\dim_{\text{H}}^\nu(X) = \dim_{\text{H}}^{(\rho_\nu)}(X).$$

2. The *packing dimension* of  $X$  with respect to  $\nu$  is

$$\text{Dim}_{\text{P}}^\nu(X) = \text{Dim}_{\text{P}}^{(\rho_\nu)}(X).$$

Note: We have assumed strong positivity here for clarity of presentation, but this assumption can be weakened in various ways for various results.

When  $\nu$  is the probability measure  $\mu$ , it is generally omitted from the terminology. Thus, the *Hausdorff dimension* of  $X$  is  $\dim_{\text{H}}(X) = \dim_{\text{H}}^\mu(X)$ , and the *packing dimension* of  $X$  is  $\text{Dim}_{\text{P}}(X) = \text{Dim}_{\text{P}}^\mu(X)$ .

It was apparently Rey [35] who first noticed that the metric  $\rho_\nu$  could be used to make Billingsley dimension a special case of Hausdorff dimension. Fernau and Staiger [14] have also investigated this notion.

We now develop gale characterizations of  $\dim_{\text{H}}^\nu$  and  $\text{Dim}_{\text{P}}^\nu$ .

**Definition.** Let  $\Sigma$  be a finite alphabet with  $|\Sigma| \geq 2$ , let  $\nu$  be a probability measure on  $\Sigma^\infty$ , and let  $s \in [0, \infty)$ .

1. A  $\nu$ - $s$ -supergale is a function  $d : \Sigma^* \rightarrow [0, \infty)$  that satisfies the condition

$$d(w)\nu(w)^s \geq \sum_{a \in \Sigma} d(wa)\nu(wa)^s \tag{A.2}$$

for all  $w \in \Sigma^*$ .

2. A  $\nu$ - $s$ -gale is a  $\nu$ - $s$ -supergale that satisfies (A.2) with equality for all  $w \in \Sigma^*$ .
3. A  $\nu$ -supermartingale is a  $\nu$ -1-supergale.
4. A  $\nu$ -martingale is a  $\nu$ -1-gale.
5. An  $s$ -supergale is a  $\mu$ - $s$ -supergale.
6. An  $s$ -gale is a  $\mu$ - $s$ -gale.
7. A supermartingale is a 1-supergale.
8. A martingale is a 1-gale.

The following observation shows how gales and supergales are affected by variation of the parameter  $s$ .

**Observation A.4** [29]. *Let  $\nu$  be a probability measure on  $\Sigma^\infty$ , let  $s, s' \in [0, \infty)$  and let  $d, d' : \Sigma^* \rightarrow [0, \infty)$ . Assume that*

$$d(w)\nu(w)^s = d'(w)\nu(w)^{s'}$$

*holds for all  $w \in \Sigma^*$ .*

1.  *$d$  is a  $\nu$ - $s$ -supergale if and only if  $d'$  is a  $\nu$ - $s'$ -supergale.*
2.  *$d$  is a  $\nu$ - $s$ -gale if and only if  $d'$  is a  $\nu$ - $s'$ -gale.*

For example, if the probability measure  $\nu$  is positive, then a function  $d : \Sigma^* \rightarrow [0, \infty)$  is a  $\nu$ - $s$ -gale if and only if the function  $d' : \Sigma^* \rightarrow [0, \infty)$  defined by  $d'(w) = \nu(w)^{s-1}d(w)$  is a  $\nu$ -martingale.

Martingales were introduced by Lévy [24] and Ville [43]. They have been used extensively by Schnorr [36, 37, 38] and others in investigations of randomness and by Lutz [26, 27] and others in the development of resource-bounded measure. Gales are a convenient generalization of martingales introduced by Lutz [28, 29] in the development of effective fractal dimensions.

The following generalization of Kraft's inequality [8] is often useful.

**Lemma A.5** [29] *Let  $d$  be a  $\nu$ - $s$ -supergale, where  $\nu$  is a probability measure on  $\Sigma^\infty$  and  $s \in [0, \infty)$ . Then, for all  $w \in \Sigma^*$  and all prefix sets  $B \subseteq \Sigma^*$ ,*

$$\sum_{u \in B} d(wu)\nu(wu)^s \leq d(w)\nu(w)^s.$$

Intuitively, a  $\nu$ - $s$ -gale  $d$  is a strategy for betting on the successive symbols in a sequence  $S \in \Sigma^\infty$ . We regard the value  $d(w)$  as the amount of money that a gambler using the strategy  $d$  will have after betting on the symbols in  $w$ , if  $w$  is a prefix of  $S$ . If  $s = 1$ , then the  $\nu$ - $s$ -gale identity,

$$d(w)\nu(w)^s = \sum_{a \in \Sigma} d(wa)\nu(wa)^s, \tag{A.3}$$

ensures that the payoffs are fair in the sense that the conditional  $\nu$ -expected value of the gambler's capital after the symbol following  $w$ , given that  $w$  has occurred, is precisely  $d(w)$ , the gambler's capital after  $w$ . If  $s < 1$ , then (A.3) says that the payoffs are less than fair. If  $s > 1$ , then (A.3) says that the payoffs are more than fair. Clearly, the smaller  $s$  is, the more hostile the betting environment is.

There are two important notions of success for a supergale.

**Definition.** Let  $d$  be a  $\nu$ - $s$ -supergale, where  $\nu$  is a probability measure on  $\Sigma^\infty$  and  $s \in [0, \infty)$ , and let  $S \in \Sigma^\infty$ .

1. We say that  $d$  *succeeds* on  $S$ , and we write  $S \in S^\infty[d]$ , if  $\limsup_{t \rightarrow \infty} d(S[0..t - 1]) = \infty$ .

2. We say that  $d$  succeeds strongly on  $S$ , and we write  $S \in S_{\text{str}}^\infty[d]$ , if  $\liminf_{t \rightarrow \infty} d(S[0..t-1]) = \infty$ .

**Notation.**

Let  $\nu$  be a probability measure on  $\Sigma^\infty$ , and let  $X \subseteq \Sigma^\infty$ .

1.  $\mathcal{G}^\nu(X)$  is the set of all  $s \in [0, \infty)$  such that there is a  $\nu$ - $s$ -gale  $d$  for which  $X \subseteq S^\infty[d]$ .
2.  $\mathcal{G}^{\nu, \text{str}}(X)$  is the set of all  $s \in [0, \infty)$  such that there is a  $\nu$ - $s$ -gale  $d$  for which  $X \subseteq S_{\text{str}}^\infty[d]$ .
3.  $\widehat{\mathcal{G}}^\nu(X)$  is the set of all  $s \in [0, \infty)$  such that there is a  $\nu$ - $s$ -supergale  $d$  for which  $X \subseteq S^\infty[d]$ .
4.  $\widehat{\mathcal{G}}^{\nu, \text{str}}(X)$  is the set of all  $s \in [0, \infty)$  such that there is a  $\nu$ - $s$ -supergale  $d$  for which  $X \subseteq S_{\text{str}}^\infty[d]$ .

The following theorem gives useful characterizations of the classical Hausdorff and packing dimensions with respect to probability measures on sequence spaces.

**Theorem A.6** (*gale characterizations of  $\dim_{\text{H}}^\nu(X)$  and  $\text{Dim}_{\text{P}}^\nu(X)$* ). *If  $\nu$  is a strongly positive probability measure on  $\Sigma^\infty$ , then, for all  $X \subseteq \Sigma^\infty$ ,*

$$\dim_{\text{H}}^\nu(X) = \inf \mathcal{G}^\nu(X) = \inf \widehat{\mathcal{G}}^\nu(X) \tag{A.4}$$

and

$$\text{Dim}_{\text{P}}^\nu(X) = \inf \mathcal{G}^{\nu, \text{str}}(X) = \inf \widehat{\mathcal{G}}^{\nu, \text{str}}(X). \tag{A.5}$$

**Proof.** In this proof we will use the following notation, for each  $w \in \Sigma^*$ ,  $\mathbf{C}_w = \{S \in \Sigma^\infty \mid w \sqsubseteq S\}$ .

Notice that for each  $S \in \Sigma^\infty$ ,  $r > 0$ , the balls  $B(S, r) = \mathbf{C}_v$ ,  $B^o(S, r) = \mathbf{C}_w$  for some  $v, w \in \Sigma^*$ . Therefore two balls  $\mathbf{C}_w, \mathbf{C}_{w'}$  are either disjoint or one contained in the other.

In order to prove (A.4) it suffices to show that for all  $s \in [0, \infty)$ ,

$$H^s(X) = 0 \iff s \in \mathcal{G}^\nu(X)$$

First, assume that  $H^s(X) = 0$ . Then  $H_1^s(X) = 0$ , which implies that for each  $r \in \mathbb{N}$ , there is a disjoint cover  $\mathcal{B} \in \mathcal{B}_1$  such that  $\sum_{B \in \mathcal{B}} \text{diam}(B)^s < 2^{-r}$ . Let  $A_r = \{w \in \Sigma^* \mid \mathbf{C}_w \in \mathcal{B}\}$ .

We define a function  $d : \Sigma^* \rightarrow [0, \infty)$  as follows. Let  $w \in \Sigma^*$ . If there exists  $v \sqsubseteq w$  such that  $v \in A_r$  then

$$d_r(w) = \left( \frac{\nu(w)}{\nu(v)} \right)^{1-s}.$$

Otherwise,

$$d_r(w) = \sum_{\substack{u, \\ wu \in A_r}} \left( \frac{\nu(wu)}{\nu(w)} \right)^s.$$

It is routine to verify that the following conditions hold for all  $r \in \mathbb{N}$ .

- (i)  $d_r$  is a  $\nu$ - $s$ -gale.
- (ii)  $d_r(\lambda) < 2^{-r}$ .
- (iii) For all  $w \in A_r$ ,  $d_r(w) = 1$ .

Let  $d = \sum_{r=0}^{\infty} 2^r d_{2^r}$ . Notice that  $d$  is a  $\nu$ - $s$ -gale. To see that  $X \subseteq S^\infty[d]$ , let  $T \in X$ , and let  $r \in \mathbb{N}$  be arbitrary. Since  $\mathcal{B}$  covers  $X$ , there exists  $w \in A_{2^r}$  such that  $w \sqsubseteq T$ . Then by (iii) above,  $d(w) \geq 2^r d_{2^r}(w) = 2^r$ . Since  $r \in \mathbb{N}$  is arbitrary, this shows that  $T \in S^\infty[d]$ , confirming that  $X \subseteq S^\infty[d]$ .

We have now shown that  $d$  is a  $\nu$ - $s$ -gale such that  $X \subseteq S^\infty[d]$ , whence  $s \in \mathcal{G}^\nu(X)$ .

Conversely, assume that  $s \in \widehat{\mathcal{G}}^\nu(X)$ . To see that  $H^s(X) = 0$ , let  $\delta > 0$ ,  $r \in \mathbb{N}$ . It suffices to show that  $H^s(X) \leq 2^{-r}$ . If  $X = \emptyset$  this is trivial, so assume that  $X \neq \emptyset$ .

Since  $s \in \widehat{\mathcal{G}}^\nu(X)$ , there is a  $\nu$ - $s$ -supergale  $d$  such that  $X \subseteq S^\infty[d]$ . Note that  $d(\lambda) > 0$  because  $X \neq \emptyset$ . Let

$$A = \{w \in \Sigma^* \mid \nu(w) < \delta, d(w) \geq 2^r d(\lambda) \text{ and } (\forall v)[v \sqsubset w \implies v \notin A]\}.$$

It is clear that  $A$  is a prefix set. It is also clear that  $\mathcal{B} = \{\mathbf{C}_w \mid w \in A\}$  is a  $\delta$ -cover of  $S^\infty[d]$ , and since  $X \subseteq S^\infty[d]$ ,  $\mathcal{B}$  is also a  $\delta$ -cover of  $X$ . By Lemma A.5 and the definition of  $A$ , we have

$$d(\lambda) \geq \sum_{w \in A} \nu(w)^s d(w) \geq 2^r d(\lambda) \sum_{w \in A} \nu(w)^s.$$

Since  $B \in \mathcal{C}_\delta(X)$  and  $d(\lambda) > 0$ , it follows that

$$H_\delta^s(X) \leq \sum_{w \in A} \nu(w)^s \leq 2^{-r}.$$

This completes the proof of (A.4).

The proof of (A.5) is based on the following three claims.

**Claim 1.** For each family  $X_i \subseteq \Sigma^\infty$ ,  $i \in \mathbb{N}$

$$\inf \mathcal{G}^{\nu, \text{str}}(\cup_i X_i) = \sup_i \inf \mathcal{G}^{\nu, \text{str}}(X_i).$$

**Claim 2.** For each  $X \subseteq \Sigma^\infty$ , if  $P_0^s(X) < \infty$  then  $\inf \mathcal{G}^{\nu, \text{str}}(X) \leq s$ .

**Claim 3.** For each  $X \subseteq \Sigma^\infty$ , if  $s > \inf \widehat{\mathcal{G}}^{\nu, \text{str}}(X)$  then  $P^s(X) = 0$ .

**Proof of Claim 1.** The  $\geq$  inequality follows from the definition of  $\mathcal{G}^{\nu, \text{str}}()$ . To prove that  $\inf \mathcal{G}^{\nu, \text{str}}(\cup_i X_i) \leq \sup_i \inf \mathcal{G}^{\nu, \text{str}}(X_i)$ , let  $s > \sup_i \inf \mathcal{G}^{\nu, \text{str}}(X_i)$ . Assume that  $X_i \neq \emptyset$  for every  $i$ , since otherwise the proof is similar taking only nonempty  $X_i$ 's. Then for each  $i \in \mathbb{N}$  there is a  $\nu$ - $s$ -gale  $d_i$  such that  $X_i \subseteq S_{\text{str}}^\infty[d_i]$ . We define a  $\nu$ - $s$ -gale  $d$  by

$$d(w) = \sum_i \frac{2^{-i}}{d_i(\lambda)} d_i(w)$$

for all  $w \in \Sigma^*$ . Then for each  $i$ , for any  $S \in X_i$ , we have

$$d(S[0..n-1]) \geq \frac{2^{-i}}{d_i(\lambda)} d_i(S[0..n-1])$$

for all  $n$ , so  $S \in S_{\text{str}}^\infty[d]$ . Therefore  $\cup_i X_i \subseteq S_{\text{str}}^\infty[d]$  and the claim follows.  $\square$

**Proof of Claim 2.** Assume that  $P_0^s(X) < \infty$ . Let  $\epsilon > 0$ . Let

$$A = \{w \mid w \in \Sigma^* \text{ and } C_w \cap X \neq \emptyset\}.$$

Notice that there is a constant  $c$  such that for every  $n$ ,  $\sum_{w \in A^n} \nu(w)^s < c$  and that for each  $T \in X$ , for every  $n$ ,  $T[0..n-1] \in A$ . For each  $n \in \mathbb{N}$  we define  $d_n : \Sigma^* \rightarrow [0, \infty)$  similarly to the first part of this proof, that is, let  $w \in \Sigma^*$ . If there exists  $v \sqsubseteq w$  such that  $v \in A^{=n}$  then

$$d_n(w) = \left( \frac{\nu(w)}{\nu(v)} \right)^{1-s}.$$

Otherwise,

$$d_n(w) = \sum_{\substack{u \\ wu \in A^n}} \left( \frac{\nu(wu)}{\nu(w)} \right)^s.$$

$d_n$  is a  $\nu$ - $s$ -gale,  $d_n(\lambda) = \sum_{u \in A^n} \nu(u)^s$  and for all  $w \in A^{=n}$ ,  $d_n(w) = 1$ .

Let  $d(w) = \sum_{n=0}^\infty \nu(w)^{-\epsilon} d_n(w)$ . Notice that  $d$  is a  $\nu$ - $(s + \epsilon)$ -gale. To see that  $X \subseteq S_{\text{str}}^\infty[d]$ , let  $T \in X$  and let  $n$  be arbitrary. Since  $T[0..n-1] \in A$ ,

$$d(T[0..n-1]) \geq \nu(T[0..n-1])^{-\epsilon} d_n(T[0..n-1]) \geq \nu(T[0..n-1])^{-\epsilon}.$$

Since  $\nu(T[0..n-1]) \xrightarrow{n \rightarrow \infty} 0$  this shows that  $T \in S_{\text{str}}^\infty[d]$ . Therefore  $X \subseteq S_{\text{str}}^\infty[d]$  and  $\inf \mathcal{G}^{\nu, \text{str}}(X) \leq s + \epsilon$  for arbitrary  $\epsilon$ , so the claim follows.  $\square$

**Proof of Claim 3.** Let  $s > t > \inf \widehat{\mathcal{G}}^{\nu, \text{str}}(X)$ . To see that  $P^s(X) = 0$ , let  $d$  be a  $\nu$ - $t$ -supergale such that  $X \subseteq S_{\text{str}}^\infty[d]$ . Let  $i \in \mathbb{N}$  and

$$X_i = \{T \mid \forall n \geq i, d(T[0..n-1]) > d(\lambda)\}.$$

Then  $X \subseteq \cup_i X_i$ . For each  $i \in \mathbb{N}$  we prove that  $P_0^s(X_i) = 0$ .

Let  $\delta_i = \min_{|w| \leq i} \nu(w)$ . Let  $\delta < \delta_i$  and  $B$  be a  $\delta$ -packing of  $X_i$ , then  $B \subseteq \{w \mid d(w) > d(\lambda)\}$  and  $\sum_{w \in B} \nu(w)^t \leq 1$ . Therefore  $P_0^t(X_i) \leq 1$  and  $P_0^s(X_i) = 0$  (since  $\sum_{w \in B} \nu(w)^s \leq \delta^{s-t} \xrightarrow{\delta \rightarrow 0} 0$ ). Therefore  $P^s(X) = 0$  and the claim follows.  $\square$

We next prove (A.5).  $\inf \mathcal{G}^{\nu, \text{str}}(X) \leq \text{Dim}_P^\nu(X)$  follows from Claims 1 and 2, and  $\text{Dim}_P^\nu(X) \leq \widehat{\mathcal{G}}^{\nu, \text{str}}(X)$  from Claim 3.  $\square$

We note that the case  $\nu = \mu$  of (A.4) was proven by Lutz [28], and the case  $\nu = \mu$  of (A.5) was proven by Athreya, Hitchcock, Lutz, and Mayordomo [1].

Guided by Theorem A.6, we now develop the constructive fractal  $\nu$ -dimensions.

**Definition.** A  $\nu$ - $s$ -supergale  $d$  is *constructive* if it is lower semicomputable, i.e., if there is an exactly computable function  $\widehat{d} : \Sigma^* \times \mathbb{N} \rightarrow \mathbb{Q}$  with the following two properties.

(i) For all  $w \in \Sigma^*$  and  $t \in \mathbb{N}$ ,  $\widehat{d}(w, t) \leq \widehat{d}(w, t+1) < d(w)$ .

(ii) For all  $w \in \Sigma^*$ ,  $\lim_{t \rightarrow \infty} \widehat{d}(w, t) = d(w)$ .

**Notation.** For each probability measure  $\nu$  on  $\Sigma^\infty$  and each set  $X \subseteq \Sigma^\infty$ , we define the sets  $\mathcal{G}_{\text{constr}}^\nu(X)$ ,  $\mathcal{G}_{\text{constr}}^{\nu, \text{str}}(X)$ ,  $\widehat{\mathcal{G}}_{\text{constr}}^\nu(X)$ , and  $\widehat{\mathcal{G}}_{\text{constr}}^{\nu, \text{str}}(X)$  exactly like the sets  $\mathcal{G}^\nu(X)$ ,  $\mathcal{G}^{\nu, \text{str}}(X)$ ,  $\widehat{\mathcal{G}}^\nu(X)$ , and  $\widehat{\mathcal{G}}^{\nu, \text{str}}(X)$ , respectively, except that the gales and supergales  $d$  are now required to be constructive.

**Definition.** Let  $\nu$  be a probability measure on  $\Sigma^\infty$ , and let  $X \subseteq \Sigma^\infty$ .

1. The *constructive  $\nu$ -dimension* of  $X$  is  $\text{cdim}^\nu(X) = \inf \widehat{\mathcal{G}}_{\text{constr}}^\nu(X)$ .
2. The *constructive strong  $\nu$ -dimension* of  $X$  is  $\text{cDim}^\nu(X) = \inf \widehat{\mathcal{G}}_{\text{constr}}^{\nu, \text{str}}(X)$ .
3. The *constructive dimension* of  $X$  is  $\text{cdim}(X) = \text{cdim}^\mu(X)$ .
4. The *constructive strong dimension* of  $X$  is  $\text{cDim}(X) = \text{cDim}^\mu(X)$ .

The fact that the “unhatted”  $\mathcal{G}$ -classes can be used in place of the “hatted”  $\widehat{\mathcal{G}}$ -classes is not as obvious in the constructive case as in the classical case. Nevertheless, Fenner [13] proved that this is the case for constructive  $\nu$ -dimension. (Hitchcock [19] proved this independently for the case  $\nu = \mu$ .) The case of strong  $\nu$ -dimension also holds with a more careful argument [1].

**Theorem A.7** (Fenner [13]). *If  $\nu$  is a strongly positive, computable probability measure on  $\Sigma^\infty$ , then, for all  $X \subseteq \Sigma^\infty$ ,*

$$\text{cdim}^\nu(X) = \inf \mathcal{G}_{\text{constr}}^\nu(X)$$

and

$$\text{cDim}^\nu(X) = \inf \mathcal{G}_{\text{constr}}^{\nu, \text{str}}(X).$$

A *correspondence principle* for an effective dimension is a theorem stating that the effective dimension coincides with its classical counterpart on sufficiently “simple” sets. The following such principle, proven by Hitchcock [20], extended a correspondence principle for computable dimension that was implicit in results of Staiger [39].

**Theorem A.8** (*correspondence principle for constructive dimension [20]*). *If  $X \subseteq \Sigma^\infty$  is any union (not necessarily effective) of computably closed, i.e.,  $\Pi_1^0$ , sets, then  $\text{cdim}(X) = \text{dim}_H(X)$ .*

We now define the constructive dimensions of individual sequences.

**Definition.** Let  $\nu$  be a probability measure on  $\Sigma^\infty$ , and let  $S \in \Sigma^\infty$ . Then the  $\nu$ -dimension of  $S$  is

$$\text{dim}^\nu(S) = \text{cdim}^\nu(\{S\}),$$

and the strong  $\nu$ -dimension of  $S$  is

$$\text{Dim}^\nu(S) = \text{cDim}^\nu(\{S\}).$$

## B Kolmogorov Complexity Characterizations

In this section we prove characterizations of constructive  $\nu$ -dimension and constructive strong  $\nu$ -dimension in terms of Kolmogorov complexity. These characterizations are used in the proof of our main theorem in section 6.

Let  $\Sigma$  be a finite alphabet, with  $|\Sigma| \geq 2$ . The *Kolmogorov complexity* of a string  $w \in \Sigma^*$  is the natural number

$$K(w) = \min \{ |\pi| \mid \pi \in \{0, 1\}^* \text{ and } U(\pi) = w \},$$

where  $U$  is a fixed optimal universal prefix Turing machine. This is a standard notion of (prefix) Kolmogorov complexity. The reader is referred to the standard text by Li and Vitanyi [25] for background on prefix Turing machines and Kolmogorov complexity.

If  $\nu$  is a probability measure on  $\Sigma^\infty$ , then the *Shannon self information* of a string  $w \in \Sigma^*$  with respect to  $\nu$  is

$$\mathcal{I}_\nu(w) = \log \frac{1}{\nu(w)}.$$

Note that  $0 \leq \mathcal{I}_\nu(w) \leq \infty$ . Equality holds on the left here if and only if  $\nu(w) = 1$ , and equality holds on the right if and only if  $\nu(w) = 0$ . Since our results here concern strongly positive probability measures, we will have  $0 < \mathcal{I}_\nu(w) < \infty$  for all  $w \in \Sigma^+$ .

The following result is the main theorem of this section. It gives characterizations of the  $\nu$ -dimensions and the strong  $\nu$ -dimensions of sequences in terms of Kolmogorov complexity.

**Theorem B.1** *If  $\nu$  is a strongly positive, computable probability measure on  $\Sigma^\infty$ , then, for all  $S \in \Sigma^\infty$ ,*

$$\dim^\nu(S) = \liminf_{m \rightarrow \infty} \frac{K(S[0..m-1])}{\mathcal{I}_\nu(S[0..m-1])} \quad (\text{B.1})$$

and

$$\text{Dim}^\nu(S) = \limsup_{m \rightarrow \infty} \frac{K(S[0..m-1])}{\mathcal{I}_\nu(S[0..m-1])}. \quad (\text{B.2})$$

**Proof.** Let  $S \in \Sigma^\infty$ . Let  $s > s' > \liminf_m \frac{K(S[0..m-1])}{\mathcal{I}_\nu(S[0..m-1])}$ . For infinitely many  $m$ ,  $K(S[0..m-1]) < s' \mathcal{I}_\nu(S[0..m-1])$ , so  $\nu(S[0..m-1])^{s'} < 2^{-K(S[0..m-1])}$ .

Let  $m \in \mathbb{N}$ . We define the computably enumerable (c.e.) set

$$A = \{ w \mid K(w) < s' \mathcal{I}_\nu(w) \},$$

and the  $\nu$ - $s$ -constructive supergale  $d_m$  as follows. If there exists  $v \sqsubseteq w$  such that  $v \in A^m$  then

$$d_m(w) = \left( \frac{\nu(w)}{\nu(v)} \right)^{1-s}.$$

Otherwise,

$$d_m(w) = \sum_{\substack{u \\ wu \in A^m}} \left( \frac{\nu(wu)}{\nu(w)} \right)^s.$$

First notice that  $d_m$  is well-defined since

$$d_m(\lambda) = \sum_{u \in A^m} \nu(u)^s \leq \sum_{u \in A^m} 2^{-K(u)} (1 - \delta)^{m(s-s')} \leq (1 - \delta)^{m(s-s')}$$

for  $\delta \in (0, 1)$  a constant testifying that  $\nu$  is strongly positive.

We define the  $\nu$ - $s$ -constructive supergale

$$d(w) = \sum_m (1 - \delta)^{-m(s-s')} d_{2m}(w) + \sum_m (1 - \delta)^{-m(s-s')} d_{2m+1}(w).$$

Notice that the fact that  $A$  is c.e. is necessary for the constructivity of  $d$ . Since for  $w \in A$ ,  $d_{|w|}(w) = 1$  we have that  $d(w) \geq (1 - \delta)^{-|w|(s-s')/2}$  for  $w \in A$ . Since for infinitely many  $m$ ,  $S[0..m-1] \in A$  we have that  $S \in S^\infty[d]$  and  $\dim^\nu(S) \leq s$ . This finishes the proof of the first inequality of (B.1).

For the other direction, let  $s > \dim^\nu(S)$ . Let  $d$  be a  $\nu$ - $s$ -constructive gale succeeding on  $S$ . Let  $c \geq d(\lambda)$  be a rational number.

Let  $B = \{w \mid d(w) > c\}$ , notice that  $B$  is c.e. For every  $m$ ,

$$\sum_{w \in B^m} \nu(w)^s \leq 1.$$

Let  $\theta_m : B^m \rightarrow \{0, 1\}^*$  be the Shannon-Fano-Elias code (see for example [8]) given by the probability submeasure  $p$  defined as  $p(w) = \nu(w)^s$  for  $w \in B^m$ . Specifically, for each  $w \in B^m$ ,  $\theta_m(w)$  is defined as the most significant  $1 + \lceil \log \frac{1}{p(w)} \rceil$  bits of the real number

$$\sum_{\substack{|v|=m, \\ v <_B w}} p(v) + \frac{1}{2} p(w)$$

where  $<_B$  corresponds to the words in  $B$  ordered according to their appearance in the computable enumeration of  $B$ .

Then

$$|\theta_m(w)| = 1 + \left\lceil \log \frac{1}{p(w)} \right\rceil = 1 + \lceil s \mathcal{I}_\nu(w) \rceil$$

for  $w \in B^m$ .

Since  $B$  is c.e. codification and decodification can be computed given the length, that is, every  $w \in B$  can be computed from  $|w|$  and  $\theta_{|w|}(w)$ . Therefore if  $w \in B$ ,  $K(w) \leq 2 + s \mathcal{I}_\nu(w) + 2 \log(|w|)$ .

Notice that since  $\nu$  is strongly positive,  $\mathcal{I}_\nu(w) = \Omega(|w|)$  and since there exist infinitely many  $m$  for which  $S[0..m-1] \in B$ ,

$$\liminf_{m \rightarrow \infty} \frac{K(S[0..m-1])}{\mathcal{I}_\nu(S[0..m-1])} \leq s.$$

The proof of (B.2) is analogous. □

If  $\nu$  is a strongly positive probability measure on  $\Sigma^\infty$ , then there is a real constant  $\alpha > 0$  such that, for all  $w \in \Sigma^*$ ,  $\mathcal{I}_\nu(w) \geq \alpha|w|$ . Since other notions of Kolmogorov

complexity, such as the plain complexity  $C(w)$  and the monotone complexity  $K_m(w)$  [25], differ from  $K(w)$  by at most  $O(\log |w|)$ , it follows that Theorem B.1 also holds with  $K(S[0..m-1])$  replaced by  $C(S[0..m-1])$ ,  $K_m(S[0..m-1])$ , etc.

The following known characterizations of dimension and strong dimension are simply the special case of Theorem B.1 in which  $\Sigma = \{0, 1\}$  and  $\nu = \mu$ .

**Corollary B.2** [32, 1] *For all  $S \in \{0, 1\}^\infty$ ,*

$$\dim(S) = \liminf_{m \rightarrow \infty} \frac{K(S[0..m-1])}{m}$$

and

$$\text{Dim}(S) = \limsup_{m \rightarrow \infty} \frac{K(S[0..m-1])}{m}.$$

Although the constructive dimensions have primarily been investigated in sequence spaces  $\Sigma^\infty$ , they work equally well in Euclidean spaces  $\mathbb{R}^n$ . One of several equivalent ways to achieve this is to fix a base  $k \geq 2$  in which to expand the coordinates of each point  $x = (x_1, \dots, x_n) \in \mathbb{R}^n$ . If the expansions of the fractional parts of these coordinates are  $S_1, \dots, S_n \in \Sigma_k^\infty$ , respectively, where  $\Sigma_k = \{0, \dots, k-1\}$ , and if  $S$  is the interleaving of these sequences, i.e.,

$$S = S_1[0]S_2[0] \dots S_n[0]S_1[1]S_2[1] \dots S_n[1]S_1[2]S_2[2] \dots,$$

then the *dimension* of the point  $x$  is

$$\dim(x) = n \dim(S), \tag{B.3}$$

and the *strong dimension* of  $x$  is

$$\text{Dim}(x) = n \text{Dim}(S). \tag{B.4}$$

We define the dimension and strong dimension of a point  $x$  in Euclidean space as in (B.3) and (B.4). It is convenient to characterize these dimensions in terms of Kolmogorov complexity of rational approximations in Euclidean space. Specifically, for each  $x \in \mathbb{R}^n$  and  $r \in \mathbb{N}$ , we define the *Kolmogorov complexity* of  $x$  at *precision*  $r$  to be the natural number

$$K_r(x) = \min \{ K(q) \mid q \in \mathbb{Q}^n \text{ and } |q - x| \leq 2^{-r} \}.$$

That is,  $K_r(x)$  is the minimum length of any program  $\pi \in \{0, 1\}^*$  for which  $U(\pi) \in \mathbb{Q}^n \cap B(x, 2^{-r})$ . (Related notions of approximate Kolmogorov complexity have recently been considered by Vitanyi and Vereshchagin [42] and Fortnow, Lee and Vereshchagin [15].) We also mention the quantity

$$K_r(r, x) = \min \{ K(r, q) \mid q \in \mathbb{Q}^n \text{ and } |q - x| \leq 2^{-r} \},$$

in which the program  $\pi$  must specify the precision parameter  $r$  as well as a rational approximation  $q$  of  $x$  to within  $2^{-r}$ . The following relationship between these two quantities is easily verified by standard techniques.

**Observation B.3** *There exist constants  $a, b \in \mathbb{N}$  such that, for all  $x \in \mathbb{R}^n$  and  $r \in \mathbb{N}$ ,*

$$K_r(x) - a \leq K_r(r, x) \leq K_r(x) + K(r) + b.$$

We now show that the quantity  $K_r(r, x)$  is within a constant of the Kolmogorov complexity of the first  $nr$  bits of an interleaved binary expansion of the fractional part of the coordinates of  $x$ , which was defined in section 1, together with the integer part of  $x$ .

**Lemma B.4** *There is a constant  $c \in \mathbb{N}$  such that, for all  $x = (x_1, \dots, x_n) \in \mathbb{R}^n$ , all interleaved binary expansions  $S$  of the fractional parts of  $x_1, \dots, x_n$ , and all  $r \in \mathbb{N}$ ,*

$$|K_r(r, x) - K(\lfloor x \rfloor, S[0..nr - 1])| \leq c. \quad (\text{B.5})$$

where  $\lfloor x \rfloor$  is the interleaved binary expansion of  $(\lfloor x_1 \rfloor, \dots, \lfloor x_n \rfloor)$

**Proof.** We first consider the case  $x \in [0, 1]^n$ . For convenience, let  $l = \lceil \frac{\log n}{2} \rceil$  (notice that both  $n$  and  $l$  are constants). Let  $M$  be a prefix Turing machine such that, if  $\pi \in \{0, 1\}^*$  is a program such that  $U(\pi) = w \in \{0, 1\}^*$  and  $|w|$  is divisible by  $n$ , and if  $v \in \{0, 1\}^{nl}$ , then  $M(\pi v) = (|w|/n, q)$ , where  $q \in \mathbb{Q}^n$  is the dyadic rational point whose interleaved binary expansion is  $wv$ . Let  $c_1 = nl + c_M$ , where  $c_M$  is an optimality constant for  $M$ . Let  $x \in \mathbb{R}^n$ , let  $S$  be an interleaved binary expansion of  $x$ , and let  $r \in \mathbb{N}$ . Let  $\pi \in \{0, 1\}^*$  be a witness to the value of  $K_r(S[0..nr - 1])$ , and let  $v = S[nr..n(r + l) - 1]$ . Then  $M(\pi v) = (r, q)$ , where  $q$  is the dyadic rational point whose interleaved binary expansion is  $S[0, n(l + r) - 1]$ . Since

$$|q - x| = \sqrt{n(2^{-(r+l)})^2} = 2^{-(r+l)}\sqrt{n} \leq 2^{-r},$$

it follows that

$$K_r(r, x) \leq |\pi v| + c_M = K(S[0..nr - 1]) + c_1, \quad (\text{B.6})$$

which is one of the inequalities we need to get (B.5).

We now turn to the reverse inequality. For each  $r \in \mathbb{N}$  and  $q \in \mathbb{Q}^n$ , let  $A_{r,q}$  be the set of all  $r$ -dyadic points within  $2^{l-r} + 2^r$  of  $q$ . That is,  $A_{r,q}$  is the set of all points  $q' = (q'_1, \dots, q'_n) \in \mathbb{Q}^n$  such that  $|q - q'| \leq 2^{l-r} + 2^r$  and each  $q'_i$  is of the form  $2^{-r}a'_i$  for some integer  $a'_i$ .

Let  $q', q'' \in A_{r,q}$ . For each  $1 \leq i \leq n$ , let  $q'_i = 2^{-r}a'_i$  and  $q''_i = 2^{-r}a''_i$  be the  $i$ th coordinates of  $q'$  and  $q''$ , respectively. Then, for each  $1 \leq i \leq n$ , we have

$$\begin{aligned} |a'_i - a''_i| &= 2^r |q'_i - q''_i| \\ &\leq 2^r (|q' - q| + |q'' - q|) \\ &\leq 2^{r+1} (2^{l-r} + 2^{-r}) \\ &= 2^{l+1} + 2. \end{aligned}$$

This shows that there are at most  $2^{l+1} + 3$  possible values of  $a'_i$ . It follows that

$$|A_{r,q}| \leq (2^{l+1} + 3)^n. \quad (\text{B.7})$$

Let  $M'$  be a prefix Turing machine such that, if  $\pi \in \{0, 1\}^*$  is a program such that  $U(\pi) = (r, q) \in \mathbb{N} \times \mathbb{Q}^n$ ,  $0 \leq m < |A_{r,q}|$ , and  $s_m$  is the  $m$ th string in the standard enumeration  $s_0, s_1, s_2, \dots$  of  $\{0, 1\}^*$ , then  $M'(\pi 0^{|s_m|} 1 s_m)$  is the  $nr$ -bit interleaved binary expansion of the fractional points of the coordinates of the  $m$ th element of a canonical enumeration of  $A_{r,q}$ . Let  $c_2 = n(2^{l'} + 1) + c_{M'}$ , where  $l' = \lceil \log(2^{l'+1} + 3) \rceil$  and  $c_{M'}$  is an optimality constant for  $M'$ .

Let  $x \in \mathbb{R}^n$ , let  $S$  be an interleaved binary expansion of  $x$ , and let  $r \in \mathbb{N}$ . Let  $q'$  be the  $r$ -dyadic point whose interleaved binary expansion is  $S[0..nr-1]$ , and let  $\pi \in \{0, 1\}^*$  be a witness to the value of  $K_r(r, x)$ . Then  $U(\pi) = (r, q)$  for some  $q \in \mathbb{Q}^n \cap B(x, 2^{-r})$ . Since

$$\begin{aligned} |q' - q| &\leq |q' - x| + |q - x| \\ &\leq 2^{-r} \sqrt{n} + 2^{-r} \\ &\leq 2^{l-r} + 2^{-r}, \end{aligned}$$

we have  $q' \in A_{q,r}$ . It follows that there exists  $0 \leq m < |A_{r,q}|$  such that  $M'(\pi 0^{|s_m|} 1 s_m) = S[0..nr-1]$ . This implies that

$$\begin{aligned} K(S[0..nr-1]) &\leq |\pi 0^{|s_m|} 1 s_m| + c_{M'} \\ &= K_r(r, x) + 2|s_m| + c_{M'} + 1 \\ &\leq K_r(r, x) + 2|s_{|A_{r,q}|-1}| + c_{M'} + 1 \\ &= K_r(r, x) + 2\lceil \log |A_{r,q}| \rceil + c_{M'} + 1 \\ &\leq K_r(r, x) + 2\lceil n \log(2^{l'+1} + 3) \rceil + c_{M'} + 1 \\ &\leq K_r(r, x) + c_2. \end{aligned} \tag{B.8}$$

If we let  $c = \max\{c_1, c_2\}$ , then (B.6) and (B.8) imply (B.5).

For the general case, notice that  $K_r(r, \lfloor x \rfloor) = K(\lfloor x \rfloor) + O(1)$ .  $\square$

We now have the following characterizations of the dimensions and strong dimensions of points in Euclidean space.

**Theorem B.5** *For all  $x \in \mathbb{R}^n$ ,*

$$\dim(x) = \liminf_{r \rightarrow \infty} \frac{K_r(x)}{r}, \tag{B.9}$$

and

$$\text{Dim}(x) = \limsup_{r \rightarrow \infty} \frac{K_r(x)}{r}. \tag{B.10}$$

**Proof.** Let  $x \in \mathbb{R}^n$ , and let  $S$  be an interleaved binary expansion of the fractional parts of the coordinates of  $x$ . By (B.3) and Corollary B.2, we have

$$\begin{aligned} \dim(x) &= n \dim(S) \\ &= n \liminf_{m \rightarrow \infty} \frac{K(S[0..m-1])}{m}. \end{aligned}$$

Since all values of  $K(S[0..m-1])$  with  $nr \leq m < n(r+1)$  are within a constant (that depends on the constant  $n$ ) of one another, it follows that

$$\begin{aligned} \dim(x) &= n \liminf_{r \rightarrow \infty} \frac{K(S[0..nr-1])}{nr} \\ &= \liminf_{r \rightarrow \infty} \frac{K(S[0..nr-1])}{r}. \end{aligned}$$

Since  $K(r) = O(\log r)$  [25], it follows by Observation B.3 and Lemma B.4 that (B.9) holds. The proof that (B.10) holds is analogous.  $\square$

## C Additional Proofs

**Proof of Observation 3.1.** Assume the hypothesis, with  $S = (S_0, \dots, S_{k-1})$  and  $\text{dom}(S) = D$ , and let  $c_0, \dots, c_{k-1}$  be contraction ratios of  $S_0, \dots, S_{k-1}$ , respectively. Let  $c = \min\{c_0, \dots, c_{k-1}\}$ , noting that  $c \in (0, 1)$ , and fix  $z \in D$ . Let

$$r = \frac{1}{1-c} \max_{0 \leq i < k} |S_i(z) - z|,$$

and let  $A = D \cap \mathcal{B}_r(z)$ . Then  $A$  is a closed subset of the compact set  $\mathcal{B}_r(z)$ , and  $z \in A$ , so  $A \in \mathcal{K}(S)$ . For all  $x \in A$  and  $0 \leq i < k$ , we have

$$\begin{aligned} |S_i(x) - z| &\leq |S_i(x) - S_i(z)| + |S_i(z) - z| \\ &= c|x - z| + |S_i(z) - z| \\ &\leq cr + (1-c)r \\ &= r, \end{aligned}$$

so each  $S_i(A) \subseteq A$ , so  $S(A) \subseteq A$ .  $\square$

Our proof of Observation 3.2 uses the *Hausdorff metric* on  $\mathcal{K}(\mathbb{R}^n)$ , which is the function  $\rho_H : \mathcal{K}(\mathbb{R}^n) \times \mathcal{K}(\mathbb{R}^n) \rightarrow [0, \infty)$  defined by

$$\rho_H(A, B) = \max\left\{\sup_{x \in A} \inf_{y \in B} |x - y|, \sup_{y \in B} \inf_{x \in A} |x - y|\right\}$$

for all  $A, B \in \mathcal{K}(\mathbb{R}^n)$ . It is easy to see that  $\rho_H$  is a *metric* on  $\mathcal{K}(\mathbb{R}^n)$ . It follows that  $\rho_H$  is a metric on  $\mathcal{K}(S)$  for every IFS  $S$ .

**Proof of Observation 3.2.** Assume the hypothesis, with  $S = (S_0, \dots, S_{k-1})$ , and let  $c_0, \dots, c_{k-1}$  be contraction ratios of  $S_0, \dots, S_{k-1}$ , respectively. The definition of  $\rho_H$  implies immediately that, for all  $E, F \in \mathcal{K}(S)$  and  $0 \leq i < k$ ,  $\rho_H(S_i(E), S_i(F)) = c_i \rho_H(E, F)$ . It follows by an easy induction that, if we let  $c = \max\{c_0, \dots, c_{k-1}\}$ , then, for all  $w \in \Sigma_k^*$ ,

$$\rho_H(S_A(w), S_B(w)) \leq c^{|w|} \rho_H(A, B). \quad (\text{C.1})$$

To see that  $S_A = S_B$ , let  $T \in \Sigma_k^\infty$ , and let  $\epsilon > 0$ . For each  $w \sqsubseteq T$ , (3.1), (3.2), and (3.3) tell us that

$$\begin{aligned} \rho_H(\{S_A(T)\}, S_A(w)) &\leq \text{diam}(S_A(w)) \\ &\leq c^{|w|} \text{diam}(A) \end{aligned} \quad (\text{C.2})$$

and

$$\begin{aligned}\rho_H(\{S_A(T)\}, S_B(w)) &\leq \text{diam}(S_B(w)) \\ &\leq c^{|w|} \text{diam}(B).\end{aligned}\tag{C.3}$$

Since  $c \in (0, 1)$ , (C.1), (C.2), and (C.3) tell us that there is a prefix  $w \sqsubseteq T$  such that

$$\begin{aligned}&\rho_H(\{S_A(T)\}, \{S_B(T)\}) \\ &\leq \rho_H(\{S_A(T)\}, S_A(w)) + \rho_H(S_A(w), S_B(w)) \\ &+ \rho_H(\{S_B(T)\}, S_B(w)) \\ &< \epsilon/3 + \epsilon/3 + \epsilon/3 \\ &= \epsilon.\end{aligned}$$

Since this holds for all  $\epsilon > 0$ , it follows that  $\rho_H(\{S_A(T)\}, \{S_B(T)\}) = 0$ , i.e., that  $S_A(T) = S_B(T)$ .  $\square$

**Proof of Corollary 4.2.** Let  $S$  be a computable IFS. Then  $F(S)$  is compact, hence closed, and is computable by Theorem 4.1, so  $F(S)$  is computably closed by Observation 2.2. It follows by the correspondence principle for constructive dimension (Theorem A.8) that  $\text{cdim}(F(S)) = \text{dim}_H(F(S))$ .  $\square$

**Proof of Lemma 4.3.** Assume the hypothesis, and let  $B$  be a ball of radius  $r$ . Let

$$\mathcal{M} = \{G \in \mathcal{G} \mid B \cap \overline{G} \neq \emptyset\},$$

and let  $m = |\mathcal{M}|$ . Let  $B'$  be a closed ball that is concentric with  $B$  and has radius  $(1 + 2b)r$ . Then  $B'$  contains  $\overline{G}$  for every  $G \in \mathcal{M}$ . Since each  $G \in \mathcal{M}$  contains a ball  $B_G$  of radius  $ar$ , and since these balls are disjoint, it follows that

$$\text{volume}(B') \geq \sum_{G \in \mathcal{M}} \text{volume}(B_G).$$

This implies that

$$[(1 + 2b)r]^n \geq m(ar)^n,$$

whence  $m \leq \left(\frac{1+2b}{a}\right)^n$ .  $\square$

**Proof of Lemma 4.4.** Assume the hypothesis. From oracle Turing machines computing

$S_0, \dots, S_{k-1}$ , it is routine to construct an oracle Turing machine  $M$  computing the function

$$\tilde{S} : \text{dom}(S) \times \Sigma_k^* \rightarrow \text{dom}(S)$$

defined by the recursion

$$\begin{aligned}\tilde{S}(x, \lambda) &= x, \\ \tilde{S}(x, iw) &= S_i(\tilde{S}(x, w))\end{aligned}$$

for all  $x \in \text{dom}(S)$ ,  $w \in \Sigma_k^*$ , and  $i \in \Sigma_k$ . Fix a rational point  $q \in G \cap \mathbb{Q}^n$ , and let  $C_q$  be the oracle that returns the value  $q$  on all queries, noting that

$$|M^{C_q}(w, r) - \tilde{S}(q, w)| \leq 2^{-r}\tag{C.4}$$

holds for all  $w \in \Sigma_k^*$  and  $r \in \mathbb{N}$ . Fix  $l \in \mathbb{Z}^+$  large enough to satisfy the following conditions.

(i)  $G$  contains the closed ball of radius  $2^{-l}$  about  $q$ .

(ii) For each  $i \in \Sigma_k$ ,  $2^{-l} \leq c_i$ , where  $c_i$  is the contraction ratio of  $S_i$ .

Then a routine induction shows that, for each  $w \in \Sigma_k^*$ ,  $S_G(w)$  contains the closed ball of radius  $2^{-l(1+|w|)}$  about  $\tilde{S}(q, w)$ . It follows by (C.4) that the function  $h : \Sigma_k^* \rightarrow \mathbb{Q}^n$  defined by

$$h(w) = M^{C_q}(w, l(1 + |w|))$$

is a hub function for  $S$  and  $G$ . It is clear that  $h$  is rational-valued and exactly computable.  $\square$

**Proof of Lemma 4.5.** Let  $S$ ,  $c_{\min}$ , and  $\alpha$  be as given, and let  $c_0, \dots, c_{k-1}$  be the contraction ratios of  $S_0, \dots, S_{k-1}$ , respectively. Let  $c_{\max} = \max\{c_0, \dots, c_{k-1}\}$ , and let  $\delta = \frac{1}{2} \min\{c_{\min}, 1 - c_{\max}\}$ , noting that  $\delta \in (0, \frac{1}{2k}]$ . Since  $S$  is computable, there is, for each  $i \in \Sigma_k$ , an exactly computable function

$$\widehat{c}_i : \mathbb{N} \rightarrow \mathbb{Q} \cap [\delta, 1 - \delta]$$

such that, for all  $t \in \mathbb{N}$ ,

$$|\widehat{c}_i(t) - c_i| \leq 2^{-t}. \quad (\text{C.5})$$

For all  $T \in \Sigma_k^\infty$  and  $l, t \in \mathbb{N}$ , we have

$$\begin{aligned} & \prod_{i=0}^{l-1} \widehat{c}_{T[i]}(t+i) - \prod_{i=0}^{l-1} c_{T[i]} \\ &= \sum_{i=0}^{l-1} \left[ \left( \prod_{j=0}^{i-1} c_{T[j]} \right) \left( \prod_{j=i}^{l-1} \widehat{c}_{T[j]}(t+j) \right) \right. \\ & \quad \left. - \left( \prod_{j=0}^i c_{T[j]} \right) \left( \prod_{j=i+1}^{l-1} \widehat{c}_{T[j]}(t+j) \right) \right] \\ &= \sum_{i=0}^{l-1} (\widehat{c}_{T[i]}(t+i) - c_{T[i]}) p_i, \end{aligned}$$

where

$$p_i = \left( \prod_{j=0}^i c_{T[j]} \right) \left( \prod_{j=i+1}^{l-1} \widehat{c}_{T[j]}(t+j) \right).$$

Since each  $|p_i| \leq 1$ , it follows by (C.5) that

$$\left| \prod_{i=0}^{l-1} \widehat{c}_{T[i]}(t+i) - \prod_{i=0}^{l-1} c_{T[i]} \right| < 2^{1-t} \quad (\text{C.6})$$

holds for all  $T \in \Sigma_k^\infty$  and  $l, t \in \mathbb{N}$ .

By (4.1), we have  $2^{-\alpha/\dim(S)}/c_{\min} < 1$ , so we can fix  $m \in \mathbb{Z}^+$  such that

$$2^{1-m} < 1 - 2^{-\alpha/\dim(S)}/c_{\min}. \quad (\text{C.7})$$

For each  $T \in \Sigma_k^\infty$  and  $r \in \mathbb{N}$ , let

$$l_r(T) = \min \left\{ l \in \mathbb{N} \mid \prod_{i=0}^{l-1} \widehat{c}_{T[i]}(r + m + i + 1) \leq 2^{-r} - 2^{-(r+m)} \right\},$$

and let

$$A = \{(r, l_r(T)) \mid T \in \Sigma_k^\infty \text{ and } r \in \mathbb{N}\}.$$

Since the functions  $\widehat{c}_0, \dots, \widehat{c}_{k-1}$  are rational-valued and exactly computable, the set  $A$  is decidable. It is clear that each  $A_r$  has properties (i) and (ii).

Let  $r \in \mathbb{N}$ . To see that  $A_r$  has property (iii), let  $w \in A_r$ . Let  $l = |w|$ , and fix  $T \in \Sigma_k^\infty$  such that  $l = l_r(T)$  and  $w = T[0..l-1]$ . By the definition of  $l_r(T)$  and (C.6), we have

$$\prod_{i=0}^{l-1} c_{w[i]} < 2^{-r},$$

which implies that

$$\mathcal{I}_S(w) > r \dim(S). \quad (\text{C.8})$$

If  $l > 0$ , then the minimality of  $l_r(T)$  tells us that

$$\prod_{i=0}^{l-2} \widehat{c}_{w[i]}(r + m + i + 1) > 2^{-r} - 2^{-(r+m)}.$$

It follows by (C.6) and (C.7) that

$$\begin{aligned} \prod_{i=0}^{l-2} c_{w[i]} &> 2^{-r} - 2^{1-(r+m)} \\ &= 2^{-r} (1 - 2^{1-m}) \\ &> 2^{-(r+\alpha/\dim(S))} / c_{\min}, \end{aligned}$$

whence

$$\begin{aligned} \prod_{i=0}^{l-1} c_{w[i]} &> \frac{c_{w[l-1]}}{c_{\min}} 2^{-(r+\alpha/\dim(S))} \\ &\geq 2^{-(r+\alpha/\dim(S))}. \end{aligned}$$

This implies that

$$\pi_S(w) > 2^{-(r \dim(S) + \alpha)}. \quad (\text{C.9})$$

If  $l = 0$ , then  $\pi_S(w) = 1$ , so (C.9) again holds. Hence, in any case, we have

$$\mathcal{I}_S(w) < r \dim(S) + \alpha. \quad (\text{C.10})$$

By (C.8) and (C.10),  $A_r$  has property (iii).  $\square$

**Proof of Corollary 4.7.** Let  $F \subseteq \mathbb{R}^n$  be self-similar. Then there is an IFS  $S$  satisfying  $F(S) = F$  and the open set condition. For any such  $S$ , there is an oracle  $A \subseteq \{0, 1\}^*$  relative to which  $S$  is computable. We then have

$$\begin{aligned}
\dim_{\mathbb{H}}(F) &\leq \dim_{\mathbb{P}}(F) \\
&= \dim_{\mathbb{P}}^A(F) \\
&\leq \text{cDim}^A(F) \\
&= \sup_{x \in F} \text{Dim}^A(x) \\
&= \text{}^{(a)} \dim(F) \sup_{T \in \Sigma_k^\infty} \text{Dim}_{\pi_S}^A(T) \\
&= \dim(F) \\
&= \dim(F) \sup_{T \in \Sigma_k^\infty} \dim_{\pi_S}^A(T) \\
&= \text{}^{(b)} \sup_{x \in F} \dim^A(x) \\
&= \text{cdim}^A(F) \\
&= \text{}^{(c)} \dim_{\mathbb{H}}^A(F) \\
&= \dim_{\mathbb{H}}(F)
\end{aligned}$$

which implies the corollary. Equalities (a) and (b) hold by Theorem 4.6, relativized to  $A$ . Equality (c) holds by Corollary 4.2, relativized to  $A$ .

□